

MMSNP and MMSNP2

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Introduction

Definition

A set of problems has a **dichotomy** if any of its problems is either P-time or NP-complete.

Theorem (Ladner, 1975)

If $P \neq NP$, then NP has no dichotomy.

$NP = ESO$

no dichotomy

$MMSNP_2$

unknown

$CSP =_p MMSNP$

dichotomy

Theorem (Fagin, 1974)

NP equals ESO.

Theorem (Feder, Vardi, 1998)

MMSNP has a dichotomy iff CSP has.

Theorem (Bulatov, Zhuk, 2017)

CSP has a dichotomy.

$NP = ESO$

no dichotomy

$MMSNP_2$

unknown

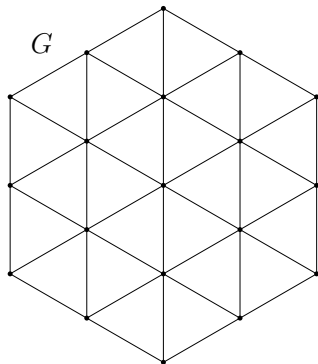
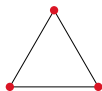
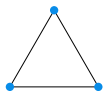
$CSP =_p MMSNP$

dichotomy

MMSNP

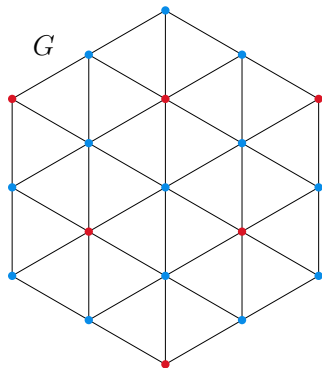
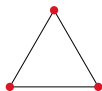
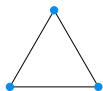
No Monochromatic Triangle

Given a graph G , colour its vertices with 2 colours so that the result omits the two following subgraphs.



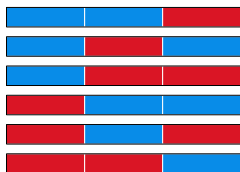
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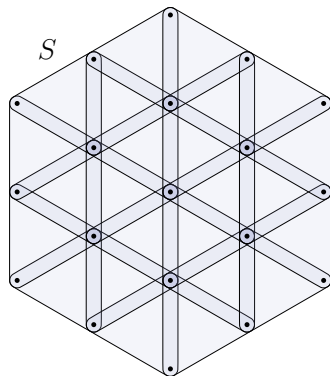


Reduction

- Replace every triangle of the input graph G with a relational triple.
- Check if the resulting structure maps to T , where T is as follows.



T



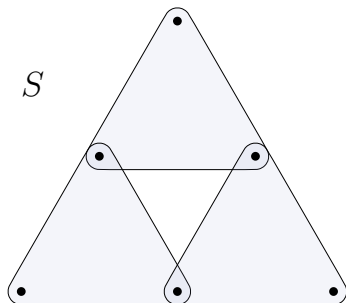
The other direction

Naive approach

- Replace every relational triple of S with a triangle.
- Check if the resulting graph G satisfies the MMSNP sentence.

Obstacle

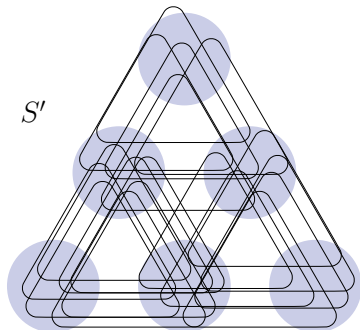
What to do when S contains implicit triangles?



Lemma (Erdős)

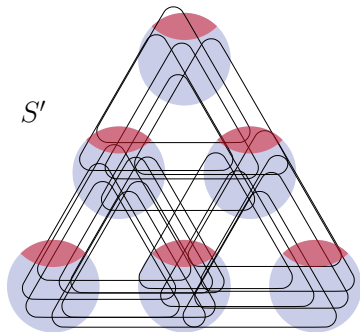
For given structures S , T , and $l > 0$ there exists S' such that

- $S \rightarrow T$ iff $S' \rightarrow T$;
- S' does not contain cycles of length less than l , i.e., the **girth** of S' is at least l .



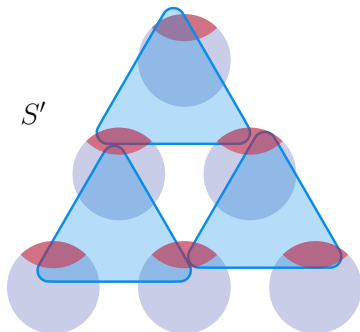
Proof

- By construction, $S' \rightarrow S$.
- The number of cycles of length $< l$ is small, so we need to remove a few tuples to get rid of them.
- If $S' \rightarrow T$, then each “bag” of size N contains at least $\frac{N}{|T|}$ vertices that are mapped to the same vertex in T .



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- If $S' \rightarrow T$, then each “bag” of size N contains at least $\frac{N}{|T|}$ vertices that are mapped to the same vertex in T .
- Tuples are distributed uniformly, so every triple of “bags” has at least one tuple induced on them.

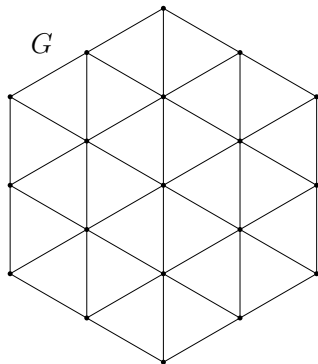
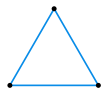


MMSNP2

Definition

No Monochromatic Edge Triangle

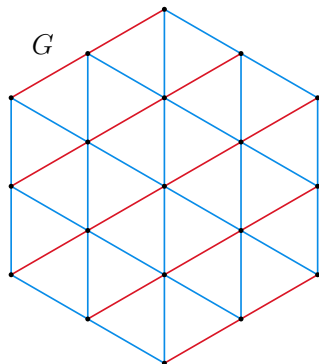
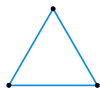
Given a graph G , colour its edges with 2 colours so that the result omits the two following subgraphs.



Definition

No Monochromatic Edge Triangle

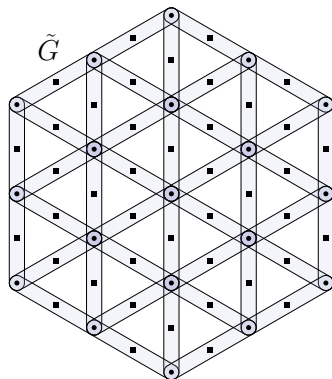
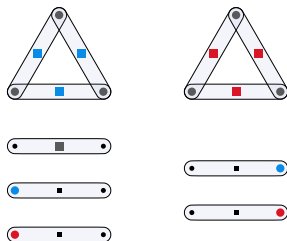
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Reduction to MMSNP

Reduction

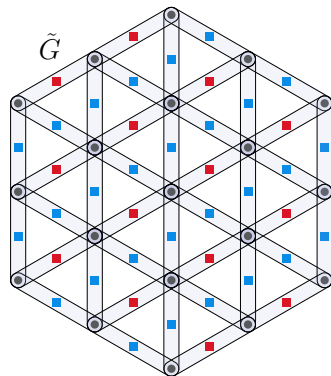
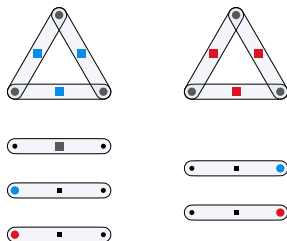
Replace every edge with a triple, where the new vertex represents the edge.



Reduction to MMSNP

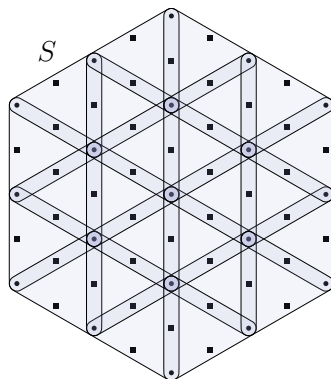
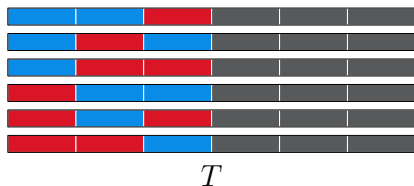
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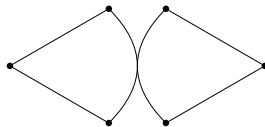
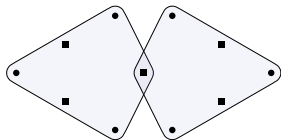
Reduction

- Replace every triangle of the input graph G with a relational 6-tuple.
- Check if the resulting structure maps to T , where T is as follows.

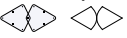


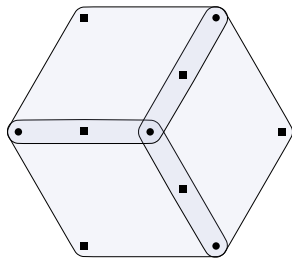
Obstacles for the other direction

- Within S' , it is not allowed to join two 6-tuples only by a vertex representing an original edge.

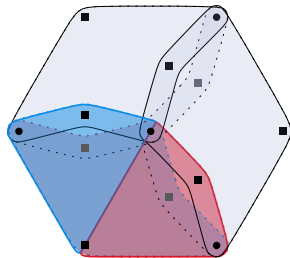


Obstacles for the other direction

- Within S' , it is not allowed to join two 6-tuples only by a vertex representing an original edge. 
- Joining 6-tuples only by vertices that represent original vertices is not sufficient.

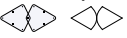


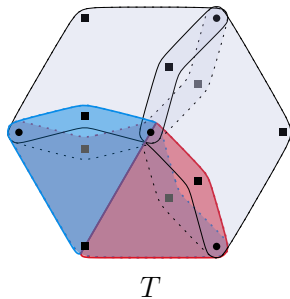
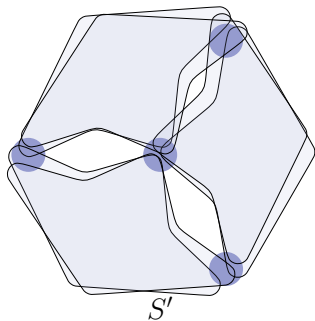
S



T

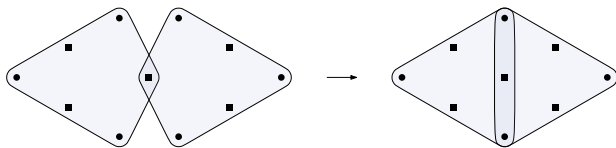
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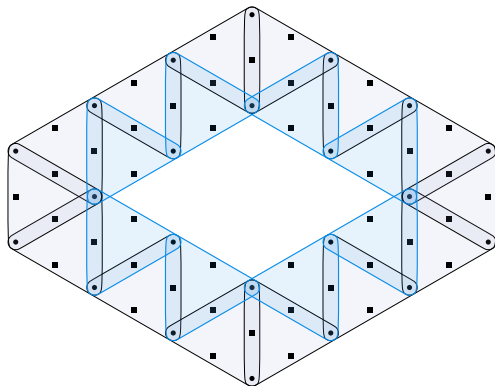
How can we help ourselves?

- We can provide that, if two tuples in S share an edge-vertex, then they share the whole implicit edge.



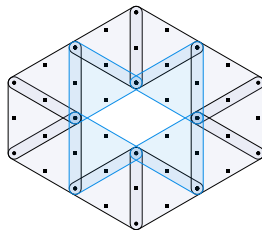
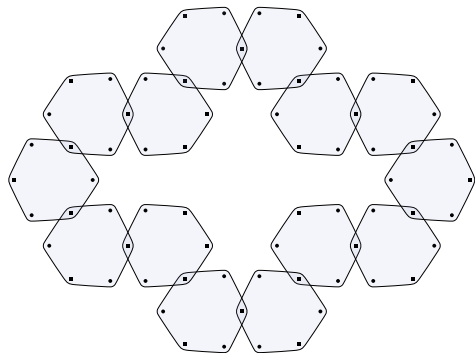
How can we help ourselves?

- We can provide that, if two tuples in S share an edge-vertex, then they share the whole implicit edge. 
- Within S' , we are allowed to join two tuples by an implicit edge, and it will not reduce the girth down to 2.



How to construct S' right?

We can apply the same Erdős' method as for MMSN_P. But then we need to identify vertices within S' in order to replace it later with a graph. This procedure reduces the girth of S' .



How to construct S' right?

- Change the measure function for vertices of S' from the Lemma of Erdős, e.g., consider the degrees of vertices.

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Thank You!