

# Reconstructing subproducts from projections

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Recall the **near unanimity (NU)** identities

$$f(y, x, x, \dots, x) \approx f(x, y, x, \dots, x) \approx \dots \approx f(x, x, \dots, x, y) \approx x$$

**NU( $l$ )**: near unanimity term of arity  $l \geq 3$

- ▶ **[Baker-Pixley]** A variety has  $\text{NU}(k + 1)$   
**iff** subproducts are determined by  $k$ -fold projections
- ▶ **[Bergman]** **If** a variety has  $\text{NU}(k + 1)$ ,  
**then** consistent systems of  $k$ -ary relations  
are  $k$ -fold projections of subproducts
- ▶ **[Our result]** A variety has  $\text{NU}(k + 2)$   
**iff** consistent systems of  $k$ -ary relations  
are  $k$ -fold projections of subproducts

**This talk:**  $k = 2$

[K. Baker, A. Pixley'75: Polynomial interpolation and the Chinese remainder theorem for algebraic systems]

## Theorem

Let  $\mathcal{V}$  be a variety. TFAE.

- (i)  $\mathcal{V}$  has NU(3).
- (ii) Every  $R \leq \mathbf{A}_1 \times \cdots \times \mathbf{A}_n$  with  $\mathbf{A}_i \in \mathcal{V}$  is uniquely determined by the system  $(\text{proj}_{ij}(R))_{i,j \in [n], i \neq j}$

### Item (ii) rephrased

- ▶ for every  $\mathbf{A}_1, \dots, \mathbf{A}_n \in \mathcal{V}$
- ▶ for every  $(P_{ij})_{i,j \in [n], i \neq j}$  where  $P_{ij} \leq \mathbf{A}_i \times \mathbf{A}_j$
- ▶ there exists **at most one**  $R \leq \mathbf{A}_1 \times \cdots \times \mathbf{A}_n$  such that  $(\forall i, j) P_{ij} = \text{proj}_{ij}(R)$

What about **at least**?

**Binary system over  $\mathcal{V}$** 

- ▶  $\mathbf{A}_1, \dots, \mathbf{A}_n \in \mathcal{V}$
- ▶  $(P_{ij})_{i,j \in [n]}$  where  $P_{ij} \leq \mathbf{A}_i \times \mathbf{A}_j$  (always  $i \neq j$ )

**Witnessing relation:**  $R \leq \mathbf{A}_1 \times \dots \times \mathbf{A}_n$  with  $(\forall i, j) P_{ij} = \text{proj}_{ij}(R)$

**Baker-Pixley:** A variety  $\mathcal{V}$  has NU(3) iff every binary system over  $\mathcal{V}$  has at most one witnessing relation

**Sometimes:** clearly no witnessing relation exists, e.g.:

- ▶  $P_{12} = \{(1, 1)\}, P_{21} = \{(1, 2)\}$
- ▶  $P_{12} = P_{23} = \{(1, 1), (2, 2)\}, P_{13} = \{(1, 2), (2, 1)\}$

**Definition**

$(P_{ij})$  is **consistent** if

- ▶  $(\forall i, j) P_{ij} = P_{ji}^{-1}$
- ▶  $(\forall i, j, k) (\forall a_i a_j \in P_{ij}) (\exists a_k) a_i a_k \in P_{ik} \text{ and } a_j a_k \in P_{jk}$

[G. Bergman'77: On the existence of subalgebras of direct products with prescribed  $d$ -fold projections]

### Theorem

*Let  $\mathcal{V}$  be a variety. Then (i) implies (ii).*

- (i)  $\mathcal{V}$  has NU(3).*
- (ii) Every consistent binary system  $(P_{ij})$  over  $\mathcal{V}$  has a witnessing relation.*

### Remarks:

- ▶ Bergman gave strengthening (ii') of (ii) and proved (i)  $\Leftrightarrow$  (ii')
- ▶ Very similar result later obtained in the context of CSPs  
[Feder, Vardi'98], [Jeavons, Cohen, Cooper'98]

[Barto, Kozik, Tan, Valeriote: Sensitive instances of CSPs, submitted]

## Theorem

Let  $\mathcal{V}$  be a variety. TFAE.

- (i)  $\mathcal{V}$  has NU(4).
- (ii) Every consistent binary system  $(P_{ij})$  over  $\mathcal{V}$  has a witnessing relation.

For a local version (concerning a single algebra):

## Definition

An algebra  $\mathbf{A}$  has **local NU( $l$ )** if for every finite  $F \subseteq A$  there exists an  $l$ -ary term operation  $t_F$  of  $\mathbf{A}$  such that  $t_F(b, a, \dots, a) = t_F(a, b, a, \dots, a) = \dots = t_F(a, \dots, a, b) = a$  for every  $a, b \in F$ .

### Theorem (Local Baker-Pixley)

Let  $\mathbf{A}$  be an idempotent algebra. TFAE.

- (i)  $\mathbf{A}$  has local NU(3).
- (ii) Every binary system over  $\{\mathbf{A}\}$  has at most one witnessing relation.

### Theorem (Local version of our result)

Let  $\mathbf{A}$  be an idempotent algebra. TFAE.

- (i)  $\mathbf{A}$  has local NU(4).
- (ii) Every binary system over  $\{\mathbf{A}^2\}$  has at least one witnessing relation.

**Remark:** Idempotency necessary, square in  $\mathbf{A}^2$  as well.

## Theorem (Local version of our result)

Let  $\mathbf{A}$  be an idempotent algebra. TFAE.

- (i)  $\mathbf{A}$  has local NU(4).
- (ii) Every binary system over  $\{\mathbf{A}^2\}$  has at least one witnessing relation.

(ii)  $\Rightarrow$  (i)

- ▶ careful choices of systems give “very local” NU(4)’s
- ▶ local NU(4)’s can be assembled from these [\[Horowitz’13\]](#)

(i)  $\Rightarrow$  (ii)

- ▶ candidate witness (the largest if any exists):  
$$R = \{a_1 a_2 \dots a_n : (\forall i, j) a_i a_j \in P_{ij}\}$$
- ▶ enough to show:  $(\forall i, j)(\forall a_i a_j \in P_{ij})$  there is an extension in  $R$



## Theorem (Local version of our result)

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(i)  $\Rightarrow$  (ii)

- ▶ predecessor: a version for finite algebras [BK]
- ▶ main tool for the predecessor: a loop lemma
- ▶ main tool for this result: an infinite loop lemma

**Here:**  $S \subseteq T \leq \mathbf{B}^2$ ,  $S$  “locally absorbs”  $T$ ,  $\mathbf{B}$  idempotent

[~ Olšák'17] If  $S$  is symmetric and  $\Delta_A \subseteq T$ , then  $S \cap \Delta_A \neq \emptyset$ .

[BKTV] If  $S$  has a directed cycle and  $\Delta_A \subseteq T$ , then  $S \cap \Delta_A \neq \emptyset$ .

[BKTV] if  $S$  has a long d.walk and  $\Delta_A \cup S^{-1} \subseteq T$ , then  $S \cap \Delta_A \neq \emptyset$ .

**Fix**  $a_1, a_2, a_3 \in B$  and **assume**

$\exists a_4$  such that  $a_i a_j \in P_{ij}$  for  $(i, j) \in \mathcal{I}$  ( $a_1 a_2 a_3 a_4$  works for  $\mathcal{I}$ )

$\exists a_4$  such that  $a_i a_j \in P_{ij}$  for  $(i, j) \in \mathcal{J}$

**Consider**

$S = \{a_4 a'_4 : a_1 a_2 a_3 a_4 \text{ works for } \mathcal{I}, a_1 a_2 a_3 a'_4 \text{ works for } \mathcal{J}\}$

$T = \{b_4 b'_4 : (\exists b_1, b_2, b_3) b_1 b_2 b_3 b_4 \text{ work for } \mathcal{I}, \dots\}$

**Then**

$S$  locally absorbs  $T$  (because of local NU)

if  $S$  and  $T$  satisfies ... we get  $a_4 a_4 \in S$  for some  $a_4$

ie.  $a_1 a_2 a_3 a_4$  works for  $\mathcal{I} \cup \mathcal{J}$

Of interest for  $\infty$ -domain CSPs:

### Question

Assume  $\mathbf{A}$  is oligomorphic core and  $\mathbf{A}$  has a *quasi-NU(4)*, i.e.,  
 $t(y, x, x, x) \approx t(x, y, x, x) \approx t(x, x, y, x) \approx t(x, x, x, y) \approx$   
 $t(x, x, x, x)$

*Does every binary system over  $\{\mathbf{A}\}$  necessarily have at least one witnessing relation?*

### Remarks:

- ▶ quasi-NU(4)  $\Rightarrow$  local NU(4), but not idempotent
- ▶ the loop lemma with quasi-absorption does not work

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**Thank you!**