# Dynamic programming in discrete time – examples

#### Martin Branda

#### Charles University in Prague Faculty of Mathematics and Physics Department of Probability and Mathematical Statistics

#### Computational Aspects of Optimization

(日)

22-05-2016

1 / 11

# Dynamic programming

Finite  ${\mathcal T}$  or infinite  $\infty$  time horizon

$$\max_{A_t,c_t} \sum_{t=1}^{T \vee \infty} \left(\frac{1}{1+i}\right)^{t-1} u(c_t)$$
  
s.t.  
$$A_t = (1+r)A_{t-1} + Y_t - c_t.$$
 (1)

- *u* utility function
- *A<sub>t</sub>* **state variables** representing total amount of resources available to the consumer.
- c<sub>t</sub> control variables maximizing the consumer's utility. It affects the resources available in the next period.
- $Y_t$  exogenous income
- 1/(1+i) discount factor, r exogenous interest rate

If we assume that there is a finite terminal period T:

$$\begin{aligned} \mathcal{V}_{1}(A_{0}) &= \max_{A_{t},c_{t}} \sum_{t=1}^{T} \left(\frac{1}{1+i}\right)^{t-1} u(c_{t}) \\ &= \max_{A_{t},c_{t}} u(c_{1}) + \frac{1}{1+i} u(c_{2}) + \dots + \left(\frac{1}{1+i}\right)^{T-1} u(c_{T}) \\ &= \max_{A_{t},c_{t}} u(c_{1}) + \frac{1}{1+i} \left[\sum_{t=2}^{T} \left(\frac{1}{1+i}\right)^{t-2} u(c_{t})\right] \\ &\text{s.t.} \\ &A_{t} = (1+r)A_{t-1} + Y_{t} - c_{t}. \end{aligned}$$

▶ ◀ ≣ ▶ ≣ ∽ � < < 22-05-2016 3 / 11

<ロ> <同> <同> < 同> < 同>

We rewrite the maximization problem recursively and obtain the **Bellman** equation

$$V_t(A_{t-1}) = \max_{A_t,c_t} u(c_t) + \frac{1}{1+i} V_{t+1}(A_t),$$

where  $A_t = (1 + r)A_{t-1} + Y_t - c_t$ .

Moreover, since u does not depend on the time period, we can write

$$V(A_{t-1}) = \max_{A_t, c_t} u(c_t) + \frac{1}{1+i} V(A_t),$$
  
= 
$$\max_{c_t} u(c_t) + \frac{1}{1+i} V((1+r)A_{t-1} + Y_t - c_t),$$

with  $V_{T+1}(A_T) = V(A_T) \equiv 0$ .

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三 ののの

## Dynamic programming

First order optimality conditions

$$\begin{array}{rcl} \frac{\partial V(A_{t-1})}{\partial c_t} & = & 0, \\ \frac{\partial V(A_{t-1})}{\partial A_{t-1}} & = & 0. \end{array}$$

In particular,

$$\begin{array}{lll} \displaystyle \frac{\partial V(A_{t-1})}{\partial c_t} & = & u'(c_t) + \frac{1}{1+i}V'(A_t)\frac{\partial A_t}{\partial c_t}, \\ \displaystyle \frac{\partial V(A_{t-1})}{\partial A_{t-1}} & = & \frac{1}{1+i}V'(A_t)\frac{\partial A_t}{\partial A_{t-1}}, \end{array}$$

where using  $A_t = (1 + r)A_{t-1} + Y_t - c_t$  we have

$$\frac{\partial A_t}{\partial c_t} = -1, \quad \frac{\partial A_t}{\partial A_{t-1}} = 1 + r.$$

イロト 不得 トイヨト イヨト 二日

#### Cake eating problem:

- $u(c) = 2c^{1/2}$ ,
- $\Pi_0 = 1, \ \Pi_T = 0,$
- $\Pi_t = \Pi_{t-1} c_t \ldots$

イロト 不得 トイヨト イヨト 二日

### Example: Cake eating problem

Bellman equation

$$V(\Pi_{t-1}) = \max_{c_t} u(c_t) + \frac{1}{1+i} V(\Pi_t),$$
  
s.t.  $\Pi_t = \Pi_{t-1} - c_t.$ 

Optimality conditions

$$\begin{array}{ll} \displaystyle \frac{\partial V(\Pi_{t-1})}{\partial c_t} & = & u'(c_t) + \frac{1}{1+i}V'(\Pi_t)\frac{\partial \Pi_t}{\partial c_t} = 0, \\ \displaystyle \frac{\partial V(\Pi_{t-1})}{\partial \Pi_{t-1}} & = & \frac{1}{1+i}V'(\Pi_t)\frac{\partial \Pi_t}{\partial \Pi_{t-1}} = 0, \end{array}$$

From  $\Pi_t = \Pi_{t-1} - c_t$ 

$$\frac{\partial \Pi_t}{\partial c_t} = -1, \ \frac{\partial \Pi_t}{\partial \Pi_{t-1}} = 1.$$
(2)

イロト 不得 トイヨト イヨト 二日

### Example: Cake eating problem

Putting them together, we obtain

$$\frac{\partial V(\Pi_{t-1})}{\partial c_t} = u'(c_t) - \frac{1}{1+i}V'(\Pi_t) = 0,$$
(3)  
$$\frac{\partial V(\Pi_{t-1})}{\partial \Pi_{t-1}} = \frac{1}{1+i}V'(\Pi_t) = 0,$$
(4)

Taking (3) for t-1

$$u'(c_{t-1}) - \frac{1}{1+i}V'(\Pi_{t-1}) = 0,$$
(5)

and plugging it into (4), we have

$$u'(c_{t-1}) = \frac{1}{1+i}u'(c_t),$$

which represents the optimal path of the cake consumption.

Martin Branda (KPMS MFF UK)

### Example: Cake eating problem

For  $u(c) = 2c^{1/2}$ , we have

$$u'(c_{t-1}) = \frac{1}{1+i}u'(c_t),$$
  

$$(c_{t-1})^{-1/2} = \frac{1}{1+i}(c_t)^{-1/2},$$
  

$$c_t = \left(\frac{1}{1+i}\right)^2 c_{t-1},$$

with initial and terminal conditions  $\Pi_0 = 1$ ,  $\Pi_T = 0$ . If we denote  $\beta = 1/(1+i)^2$ , we obtain

$$c_t = \beta c_{t-1} = \beta^{t-1} c_1.$$

Using

$$c_t = \beta c_{t-1} = \beta^{t-1} c_1.$$

and

$$\Pi_0-c_1-c_2-\ldots-c_T=\Pi_T=0,$$

we have

$$(1-\beta-\ldots-\beta^{T-1})c_1=\Pi_0,$$

and finally optimal consumption

$$\hat{c}_1 = \frac{1-\beta}{1-\beta^T} \Pi_0,$$

$$\hat{c}_t = \beta \hat{c}_{t-1} = \beta^{t-1} \hat{c}_1.$$

 M. C. Sunny Wong: Dynamic Optimization: An Introduction, Lecture Notes – University of Houston, 2013.

э

11 / 11

22-05-2016