

Finitely based 2-nilpotent Mal'tsev algebras

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Maltsev algebras and finite eq. bases

A = (A, f₁, ..., f_n) ... finite algebra ($|A| < \infty$)

A is Maltsev if \exists term m: $m(yxx) \approx m(xxy) \approx y$

Ex. rings $x-y+z$, groups $x^{-1}z$, loops $(x/y)z$, BA, ...

Id(A) ... identities holding in A

A is finitely based if $\exists \Sigma \subseteq \text{Id}(A)$: $|\Sigma| < \infty$, $\Sigma \models \text{Id}(A)$

Ex. \mathbb{Z}_n is finitely based:

$\text{Id}(\mathbb{Z}_n, +, 0, -) \models \{\text{Abelian group axioms}, n \cdot x \approx 0\}$

Finitely based Maltsev algebras

\exists non-finitely based Maltsev A ($(G, \cdot, 1, ^{-1}, c)$ '82 Bryant)

→ Which Maltsev algebras A are finitely based?

rewriting approach

↪ modulo some Σ every term has a normal form.

- affine A (e.g. \mathbb{Z}_n)
- groups (Oates, Powell '64)
- rings (L'vov/Krusc '73)
- supernilpotent A
(Vaughan-Lee '83)

HSP(A) is residually finite

↪ non-constructive

↪ also if HSP(A) has difference term (KSW '16)

←? ^{if} nilpotent A

2-nilpotent Maltsev algebras

- A is affine if $\text{Clo}(A, (\alpha)_{\text{aff}}) = \{ f \in \sum \alpha_i \cdot x_i \}$ W.R.t. a module
- A is 2-nilpotent $\Leftrightarrow \exists L, U$ affine [Freese, McKenzie '87]

$$A \simeq L \oplus U \quad A = L \times U, \quad \forall t \in \text{Clo}(A):$$

$$t^A((l_1, u_1), \dots, (l_n, u_n)) = (\underbrace{t^L(\bar{l}), t^U(\bar{u})}_{\text{affine part}}) + (\hat{t}(\bar{u}), 0)$$

, distortion'

In particular:

$$x * y := m(x, 0, y) = x + y + \hat{t}(x_u, y_u) \text{ is loop}$$

$$\hat{t}: U^n \rightarrow L$$

2-nilpotent Maltese algebras

A is 2-nilpotent

$$A \cong L \otimes U$$

$$t^A((l_1, u_1), \dots, (l_n, u_n)) = (\underbrace{t^L(\bar{l}), t^U(\bar{u})}_{\text{affine part}}) + (\hat{t}(\bar{u}), 0)$$

$\hat{t}: U^n \rightarrow L$, distortion'

rewriting approach (idea)

- ? 1) Separate affine and distortion part.
- ✓ 2) $\sum_{\text{aff.}} t =$ normal forms of affine terms
- ? 3) $\hat{\Sigma}$, to describe normal forms in the domain
 $\text{Clo } L \circ \{ \hat{t} : U^n \rightarrow L \} \circ \text{Clo } U$

2-nilpotent loops of size $p \cdot q$

$p \neq q$... primes, A nilpotent loop $|A| = p \cdot q$

Then wlog. $A \cong \mathbb{Z}_q \otimes \mathbb{Z}_p = (\mathbb{Z}_q \times \mathbb{Z}_p, *, 0, 1, \backslash)$

Lemma (\rightarrow Peter's talk)

A is term equivalent to
 $(\mathbb{Z}_q \times \mathbb{Z}_p, +, 0, -, \hat{r})$

$$\hat{r}(x, y) := (x * y)^p / (x^p * y^p)$$

$$x + y := (x * y) / r(x, y)$$

distortion $\mathbb{Z}_p^2 \rightarrow \mathbb{Z}_q$

Abelian group

$(\mathbb{Z}_p, \mathbb{Z}_q)$ - domoids

• $\sum_j \lambda_j \hat{f}(\sum_{i=1}^n \alpha_i x_i, \sum_{i=1}^n \beta_i x_i)$ form a $(\mathbb{Z}_p, \mathbb{Z}_q)$ -domoid C

Theorem (Fioravanti '19) → Patrick's talk

- 1) C is generated by unaries $\mathcal{C}^{(1)} = \{ \hat{f}(x) \in \mathcal{C} \} \leq \mathbb{Z}_q^{Z_p}$
- 2) $\hat{f} \in \mathcal{C} \iff \hat{f}(\alpha_1 x, \alpha_2 x, \dots, \alpha_n x) \in \mathcal{C}^{(1)} \forall \alpha \in \mathbb{Z}^n$
- 3) Already one $\hat{g} \in \mathcal{C}^{(1)}$ generates C

↳ A is term equivalent to $(\mathbb{Z}_q \times \mathbb{Z}_p, +, 0, -, \hat{g})$

→ full classification of Z-nl. A, $|A| = pq$, with constant.

$(\mathbb{Z}_p, \mathbb{Z}_q)$ - domoids

For $\bar{a} = (0, \dots, 0, 1, a_1, \dots, a_n) \in \mathbb{Z}_p^n$, $f \in \mathcal{C}^{(n)}$, let $\underline{f}^{\bar{a}}(\bar{x}) := \begin{cases} f(x) & \text{if } \bar{x} = (a_1, \dots, a_n) \\ 0 & \text{else} \end{cases}$

Corollary (Fioravanti '19)

If B is basis of $\mathcal{C}^{(1)}$, then

- 1) $B_1^n = \{ f^{\bar{a}} \mid f \in B \} \quad \left\{ \text{is a basis of } \mathcal{C}^{(n)} = \{ f : \mathbb{Z}_p^n \rightarrow \mathbb{Z}_q \mid f \in \mathcal{C} \} \right\}$
- 2) $B_2^n = \{ f f(\bar{a} \cdot \bar{x}) \mid f \in B \}$.

Subpower membership of A is in P: (Peter's talk)

$$Sg_{A^n}(\bar{a}_1, \dots, \bar{a}_k) = \sum_{B \in \binom{A}{1}} Sg_{B^n}(\bar{a}_1, \dots, \bar{a}_k) + \underbrace{Sg_{B_1^n}(\bar{a}_1, \dots, \bar{a}_k)}_{\text{can be computed in } |B|^n - 1}$$

A finite basis

Summary $\underline{A} = (\mathbb{Z}_q \times \mathbb{Z}_p, +, 0, -, \underbrace{\hat{f}_1, \dots, \hat{f}_n}_{B \text{ basis of } \mathcal{C}'})$

$t \in \text{Col}(A)$ has a normal form

$$t(x_1, \dots, x_n) = \underbrace{\sum_i \beta_i x_i}_{\text{affine}} + \underbrace{\sum_{\ell \in B, \bar{a}} \lambda_{\ell, \bar{a}} \hat{f}_\ell(\bar{x} \cdot \bar{x})}_{\substack{\text{lin. com.} \\ \text{over } \mathbb{Z}_2^n}}$$

Rewriting using

• $+$ is Abelian group, $(pq) \cdot x \approx 0$ Eapp

• $q \cdot \hat{f}_i(x) \approx 0$, $\hat{f}_i(x+p \cdot y) \approx \hat{f}_i(x)$, $\hat{f}_i(x + \hat{f}_i(y)) \approx \hat{f}_i(x)$

• $\hat{f}_i(c \cdot x) \approx \sum_{\ell \in B} \lambda_{\ell, i} f_\ell(x)$ for all $c \in \mathbb{Z}_p$

A is
finitely
based

What I was hoping to tell you...

For general $\underline{A} = (\mathbb{Z}_q \times \mathbb{Z}_p, m(x, y, z), f_1, \dots, f_n)$

rewriting approach

1) Separate affine and distortion part.

✓ 2) $\sum_{\text{aff.}} t$ = normal forms of affine term

✓ 3) $\hat{\Sigma}$ to describe normal forms in the clonoid

$$C = \text{Col}(\mathbb{Z}_q, +) \circ \{ \hat{t} : U \rightarrow L \} \circ \text{Col}(\mathbb{Z}_p, \underline{x-y+z})$$

$x-y+z \in \text{Col } \underline{A}$? \rightarrow Peter's talk
we don't know!!

↳ C generated by $C^{(z)}$

The clonoid part

Let \mathcal{C} be a $((\mathbb{Z}_p, \times - y + z), (\mathbb{Z}_q, +))$ -clonoid
* s.t. $\hat{f}(x \dots x) = 0 \quad \forall \hat{f} \in \mathcal{C}$

Theorem (Wymiar...?)

- 1) \mathcal{C} is generated by binary $\mathcal{C}^{(2)} \leq \mathbb{Z}_q^{\mathbb{Z}_p^2}$
 - 2) $\hat{f} \in \mathcal{C} \Leftrightarrow \hat{f}(\underbrace{\alpha_1 x + (1-\alpha_1)y, \dots, \alpha_n x + (1-\alpha_n)y}_{\text{Z-dim subspaces } \leq \mathbb{Z}_p^n \text{ containing } (xx\dots x)}) \in \mathcal{C}^{(2)}$
- If B ... basis of $\mathcal{C}^{(2)}$, then
- 3) $B_1^{(n)} = \{ f \hat{f}^{\bar{\alpha}} \mid f \in B \}$ is basis of $\mathcal{C}^{(n)}$
 - 4) $B_2^{(n)} = \{ f(x_1, \bar{\alpha} \cdot \bar{x}) \mid f \in B \}$

Do we actually need $x-y+z$?

for $\underline{A} = (\mathbb{Z}_q \times \mathbb{Z}_p, m(xyz), f_1 \dots f_n)$

for $t, s \in \text{Clo } \underline{A}$ define $t \sim s : \Leftrightarrow t_{\text{aff}} = s_{\text{aff}}$

$\mathcal{C} := \{t - s \mid t \sim s\}$ is $((\mathbb{Z}_p, \times-y+z), (\mathbb{Z}_q, +))$ -clanoid

obs.: If $f \in \text{Clo } \underline{A}, \hat{r} \in \mathcal{C} \Rightarrow f + \hat{r} \in \mathcal{C}$

Rewriting strategy:

- ~.) Find normal forms $t = t_m + \hat{t}$, $t_m \in \text{Clo}(A, m)$, $\hat{t} \in \mathcal{C}$
not affine
- ✓.) Finite identities computing \hat{t}
- ~.) Finite identities for (e.g. $m(xy) \approx m(zyx) + \hat{r}(xyz)$)
computing t_m $\hat{r} \in \mathcal{C}$

Thank you!

Any questions?