

LECTURE 7

→ A_j vari's thm, part 2

RECALL:

• $0, 1, ?$ (binary)

• Fl ... depth $\leq d$ subsystem

• From PHPD_n: UNSAT SET OF CLAUSES

in ATOT Pig.

$i \in [n+1], j \in [n]$

AIR: AJTAI'S THD (W/IMPROVEMENTS)

Удгг ЭУдо Ун>>>

in ANY FD-REENT. OF TONK PHP_n

TRUST OCCUR $\geq 2^n$ DIFHE REENT

FLAS AS SURFLAS

\Rightarrow SIZE $\geq 2^n$ σ

LAST TIME:

CLOSED UNDER SUBSETS

k -eval. of P ←

(H, S) : $\varphi \in P \longrightarrow H_\varphi \subseteq S_\varphi$: k -tree

q is (H, S) -true (OR "TRUE" w.r.t. (H, S))



$$H_q = S_q$$

LEMMA 6

ALL CLAUSES of $70h$ to PHD_n are (H, S) -~~F~~RULE if $k \leq h-2$.

LEMMA 7

IF ALL HYPOTHESES OF AN F -RULE ARE (H, S) -TRIPLE, SO IS THE CONCLUSION,

ASSUMING $k \leq \frac{h}{c_F}$

↗

CONST. DEPENDING ON F ONLY

TODAY:

$0 \leq j < 5$

"LEONORA": IF $17 \leq 2^u$ (SPACE) THEN A

k -EVAL. (M, S) EXISTS WITH $k \leq u$

appears all $q \in \Gamma$

Γ closed under suffix

PROOF THAT: $\Gamma :=$ all PLAS occurring in A
DEFAULT.

TAKE (M, S) FROM, ANSWER LG & LR

TO CONCLUDE THAT \cup IS (M, S) -TRUE.

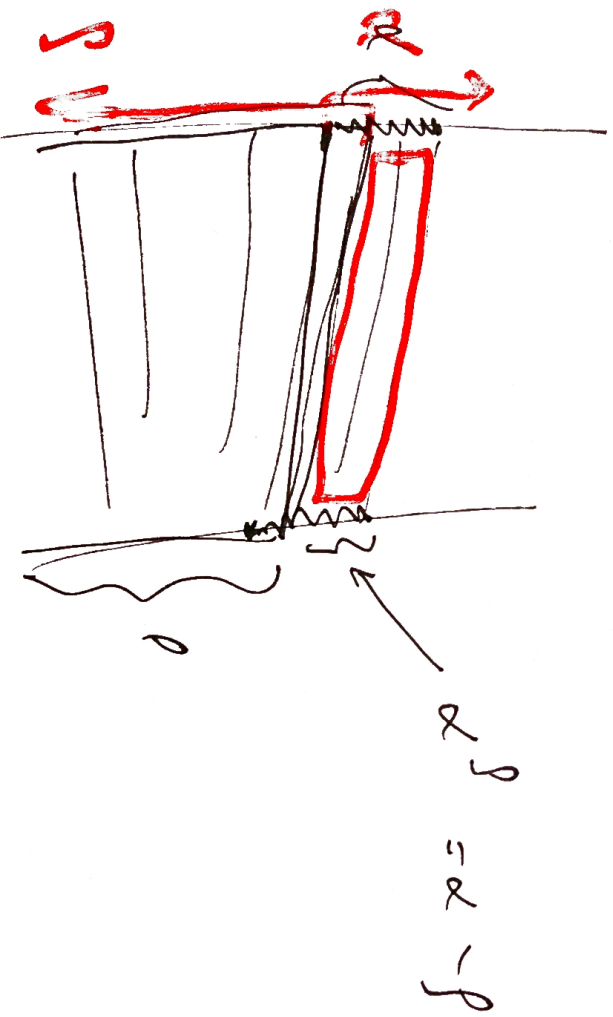
THE CAST FLA

⚡
G.

$\alpha, \beta \in \text{Draps}$ (Call Greek Letters $\in \text{Draps}$)

$\alpha \beta : \equiv \alpha - \beta$, if $\alpha \perp \beta$

β UNDEFINED, if $\alpha \perp \beta$

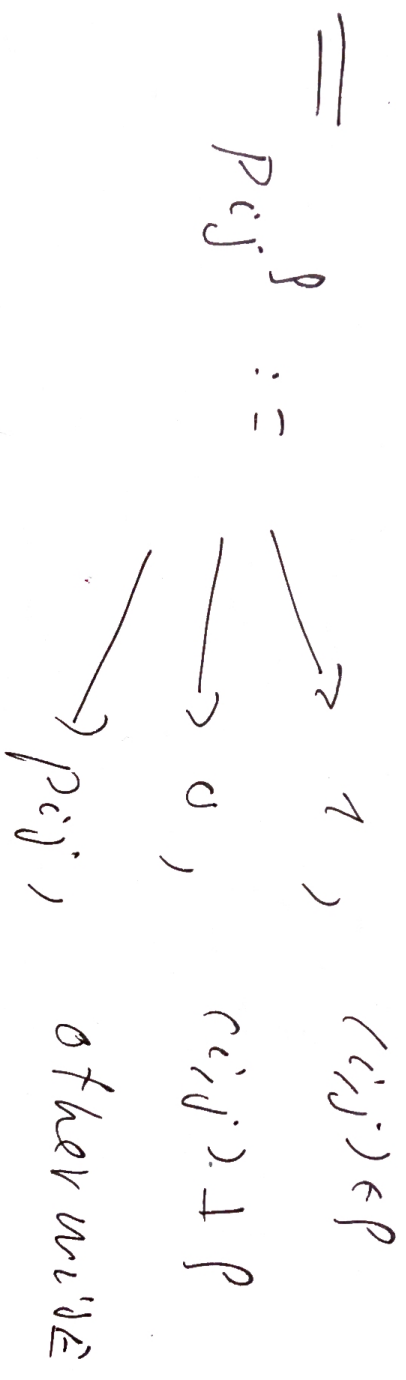


$$I - P := \chi \alpha^S \quad (\alpha \in H)$$

$$D^S := D \setminus \text{dom}(S)$$

$$R^S := R \setminus \text{rng}(S)$$

$$n_S := |R^S| \quad (\in n - |S|)$$



$$(T \text{ onto } P_H P_n)^S \cong T \text{ onto } P_H P_n^S$$

L15.2.1 Type \mathcal{O} maps

$$(i) \quad H \triangleleft S \Rightarrow H^{\mathcal{P}} \triangleleft S^{\mathcal{P}}$$

$$(ii) \quad |g| + h(S) \leq n \Rightarrow S^{\mathcal{P}} \text{ is a free } \mathcal{O}_R^{\mathcal{P}} \text{ module}$$

$$(iii) \quad H \triangleleft S \Rightarrow S^{\mathcal{P}}(H^{\mathcal{P}}) = (S(H))^{\mathcal{P}}$$

PBFS: DIRECTLY FROM DEFINITIONS. \square

[Sec. 15.2]

COROLLARY: $1 \leq k \leq n$ and (H, S) is a

k -eval. or Δ



(H, S) is a k -EVAL. or Δ^S .

PRF: (EX) A coord. eval $\varphi = \bigcup_j x_j$ is

$$\bigcup_j H_{x_j} \Rightarrow S_\varphi \quad \text{or} \quad H|_\varphi = S_\varphi \left(\bigcup_j H_{x_j} \right)$$



$$H|_\varphi = \left(\dots \right)^S$$

$$\bigcup_j H_{x_j}^S \Rightarrow S_\varphi^S = S_\varphi^S \left(\bigcup_j H_{x_j}^S \right) \quad \square$$

THE IDEA OF A CONSTRUCTION:

$P_i := \{ \varphi \in \mathcal{P} \mid \text{dp}(\varphi) \leq i \}$, $\text{sc } P_{\mathcal{P}} = 17$

(E_{k, S_0}) for $\mathcal{P}_0 \dots$ CANONICAL (JUST ATOMS = CARDS)

(H_{i, S_i}) FOR \mathcal{P}_i



$(H_{i+1, S_{i+1}})$ FOR \mathcal{P}_{i+1}

DISJUNCTIONS $\varphi \in \mathcal{P}_{i+1}, \varphi = \bigvee_j \varphi_j, \varphi_j \in \mathcal{P}_i$.

→ CANONICAL FOR $\gamma \varphi \in \mathcal{P}_{i+1}, \varphi \in \mathcal{P}_i$.

NEED



determined by (H_{i, S_i})

SO THE TASK FACING US IS:

given H FOR DRAW H_i , ONE FOR EACH DISJUNCTION $\in \mathcal{P}_{k_1, \dots, k_r}$

for k -th $S \Rightarrow H$.

WILL DO: WE FIND SUCH S **AFTER**

A RESTRICTION BY SCORE β ,

i.e. $S \Rightarrow$ **$H \beta$**

L 15.2.2 [TECHNICAL HIGH POINT]

LET $0 < \delta < \varepsilon < 1/5$, $H_i \subseteq \text{Drops}$ for $i \leq 5$.

ASSUME $h(H_i) = \max\{|x| : x \in H_i\} \leq 4$, all $i \leq 5$,

AND $n > 1$, $k \leq n^\delta$, $s \leq 2^{n^\delta}$.

THEN $\exists p \in \text{Drops}$, $n_p = n - 1/p) = n^\varepsilon$

AND

$\exists S_i \subseteq [1, s]$, S_i a k -TREE OVER D_i^j, D_i^p

(i) $S_i \cap H_i = \emptyset$, all $i \leq 5$.

(ii) $h(S_i) \leq k$, all $i \leq 5$.

□

FIN 57

L 15.2.3:

LET $0 < \delta < \epsilon < 5^{-n}$. THEN FOR ALL $n \geq 1$,

ALL P : CLOSED UNDER SUBSETS

$$\text{if } (P) \leq \delta$$

$$|P| \leq 2^{2^n \delta}$$

EXIST β , $|P| = n - n^\beta$ AND

\exists k -eval. of P^β , with $k \leq n^\delta$.

PROOF FROM L 15.2.2:

- AS OUTLINED ON SLIDE 71
- START WITH P_0 - CANONICAL

~~GIVE FOR THE STATE:~~

PICK $s_0 < 1/5 : s = s^d$.

PART: $p_0 := \phi$, $(1/10, s_0) \rightarrow$ THE CALCULATIONS

$i \rightarrow i+1$: WE HAVE $p_0 \leq p_1 \leq \dots \leq p_i$.

$$n_{p_i} = n - 1_{p_i} = n_{s_i}$$

AND (H_{i, s_i}) : n^J -EVAL OF P_i^{j, p_i} .

APPLY L15.2.3: \dots ~~BY~~ RESTR. p OR D^{j_i}, R^{p_i} .

IS THE SAME THING AS RESTR. $p \geq p_{i+1}$.

\rightarrow GIVES US $j_i^{p_i}$

$\hookrightarrow (H_{i+1, s_{i+1}})$: n^J -EVAL OF $P_{i+1}^{j_i, p_{i+1}}$

NOTE: $(P_i^{j_i, p_i})^j = P_i^{j_i, p_{i+1}}$

- By 15.2.1: ALL (H, S) is a ADE n^d -EVAL'S OF P_1, P_2 .

□ 15.2.3

THIS YIELDS A: PROOF OF AJTAI'S THM

↳ SEE SCHEME 6

AFTER ρ : $(\text{Take } P_H P_n)^\rho = \text{Take } P_H P_{n\rho}$

So (H, S) n^d -EVAL. OF P_1, P_2

GIVES (H, S) n^d -EVAL OF ALL EVALS IN

A ~~PRO~~ RECUR. OF

AS $k = n^d < n^s = 1n\rho$, & ρ APPLYS.

□ THM. 16.

THE REST OF THE LECTURE:

A PROOF OF L.N. 2.2

• WE HAVE: $0 < \sigma < \varepsilon < 1/5$

$$k \leq n^\sigma, \quad s \leq 2^{n^\sigma}$$

$$\{H_{i \leq s}\}, \quad h(CH_{i \leq s}) \leq k$$

• WE WANT: $p, |P| = n - n^\varepsilon$

Ass \Rightarrow k -free S_i over D^p, R^p

$$S_i \triangleright H_{i \leq s}$$

LET US START WITH ONE A

(GENERALIZATION TO WHICH WILL
BE EASIER)

Enumerate H : h^1, h^2, \dots

[USING LETTERS h
TO BE IN LINE W
THE BOOK]

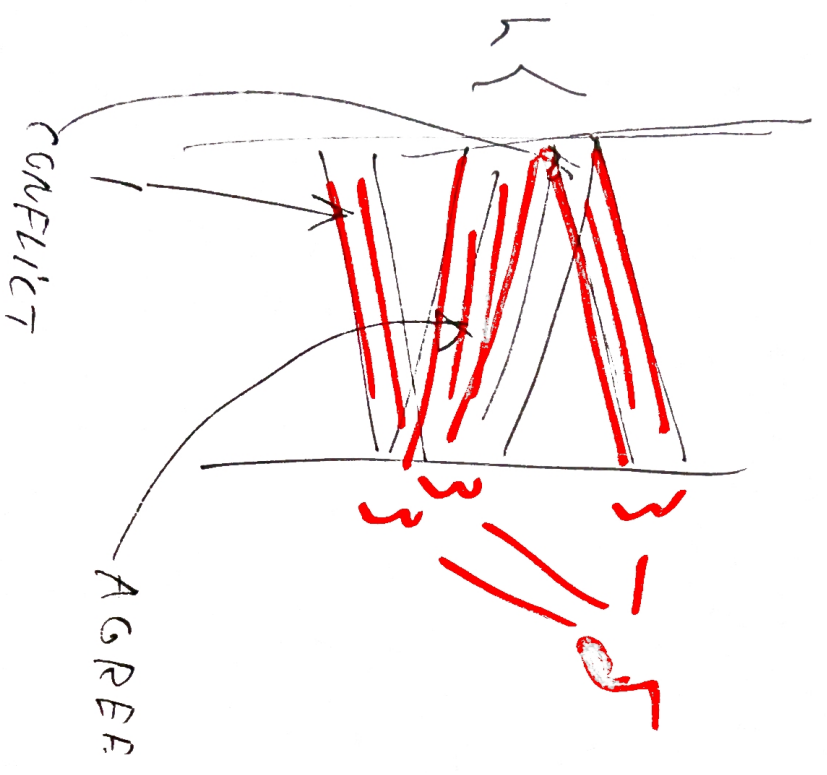
2-PLAYER GAME

- I: WILL BE CHOOSING $h \in H$
- II: WILL REPLY WITH $J \in J_{h, h}$

STEP 1:

I: picks $h_1 :=$ the FIRST h ~~set~~

II: picks $d_1 :=$ S -minimal s.t. $d_{\text{out}}(h_1) = d_{\text{in}}(d_1)$
 $\text{In}_g(h_1) = \text{Out}(d_1)$



$d_1 = h_1$ as $d_2 \geq h$ as $h \in A$

\Rightarrow GATHER STUFFS

$d_1 \perp h_1 \Rightarrow$ NEXT ROUND.

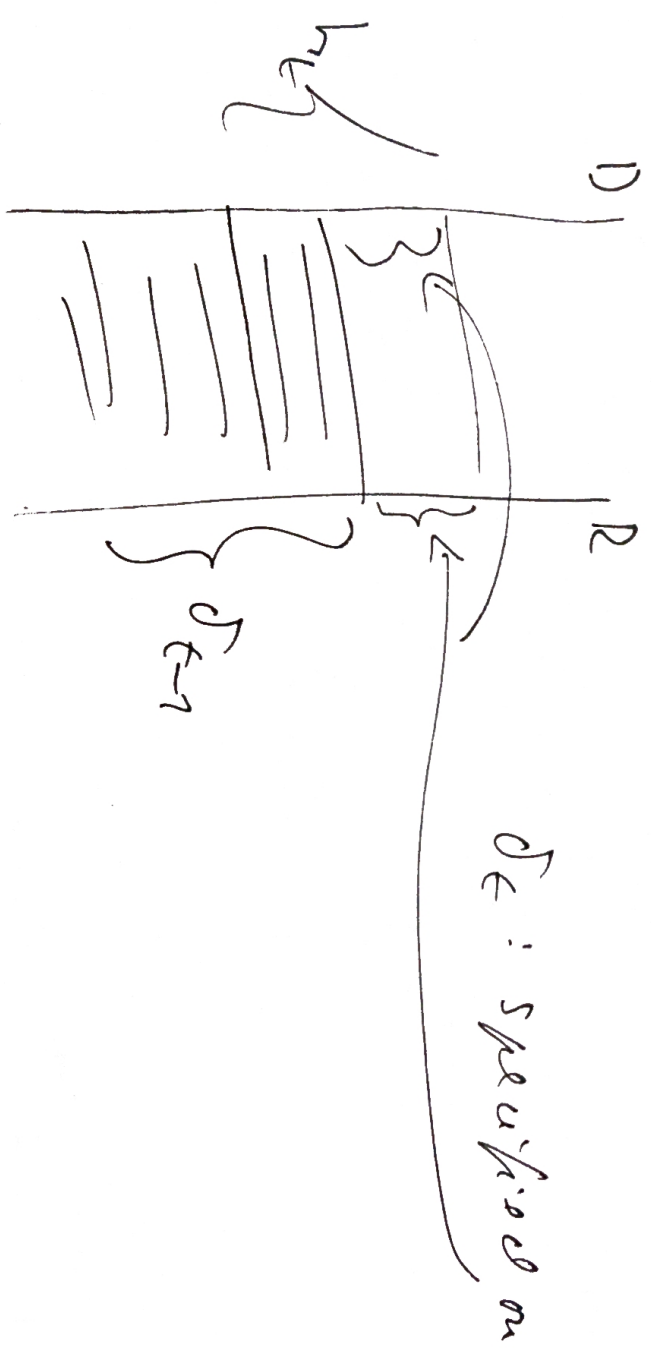
BEFORE ROUND $t \geq 2$

I: h_1, \dots, h_{t-1}

II: $\mathcal{J}_1 \subseteq \mathcal{J}_2 \subseteq \dots \subseteq \mathcal{J}_{t-1}$

ROUND t : I: picks h_t := THE FIRST $h \in H$: $h \perp \mathcal{J}_{t-1}$

II: picks $\mathcal{J}_t \stackrel{\text{def}}{=} \mathcal{J}_{t-1} \cup \text{work s.t. } \text{dom}(h_t) \subseteq \text{dom}(\mathcal{J}_t)$
 $\text{rang}(h_t) \subseteq \text{rang}(\mathcal{J}_t)$



RULES : $d_e \geq h$, so $h \in H \Rightarrow$ GAME STOPS

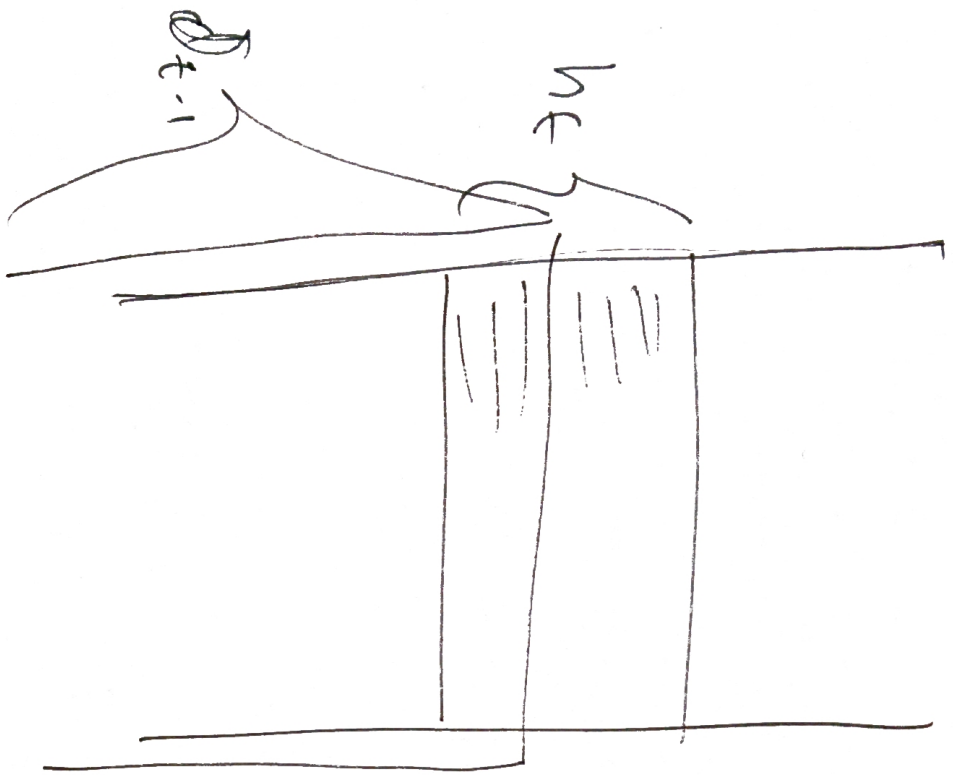
$d_e + h$, all $h \in H \Rightarrow$

(i.e. all H is EXHAUSTED)

CLAIM : $S = \{d_e \mid d_e \leq \dots \leq d_e\}$ is a FINISHED PATH
is a TREE AND $S \Rightarrow H$.

$$h(S) = \max\{|d_e| \mid d_e \in S\}$$

□



3

CRITICAL PATHS



OF NEEDS TO

BE SPECIFIED

HEAD.

EX: $A = \{(i, j) \mid \text{all } i, j\}$

\Rightarrow ALL ~~GOOD~~ MUST EXHAUST ALL j

$\Rightarrow \geq n/2$ CRITICAL PAIRS.

CLAIM 2: $\exists p \in \text{Drops}$, $n_p = n^2$ s.t.

Play on H^p has $\leq n/2$ CRITICAL PAIRS
 \square

\Leftarrow
 $\exists \epsilon \in S \Rightarrow 1/\delta \leq 2 - \delta/2 = 4$, i.e.
 S is A 4-TREE.

PROOF OF CLAIM 2 - 134 CONTRA DICTION

ASSUME :

$\forall p \dots \exists p_{log} \sigma_{\leq} \dots \sigma_{\leq} \sigma_{\leq}$ with $> 1/2$ CRIT. PAIRS

\Downarrow
TRUNCATE ~~σ_{\leq}~~ TO $1/2$ PAIRS
IN TOTAL

WE ARE GOING TO USE THIS

TO SPECIFY ANY p BY LESS BITS THAN

IT IS POSSIBLE BY COUNTING

∴

of p 's $>$ # of DATA SETS

PAT

$\tau := \rho \cup$ all critical pairs

.. ($\tau \in \text{Drops}$)

$$|\tau| = |\rho| + \frac{1}{2} \geq n - n^{\epsilon}$$

FROM τ WE CANNOT RECONSTRUCT ρ

BUT WE CAN FIND h_1 : THE FIRST $k \in H$, $h \parallel \tau$

- \mathcal{J}_1, \dots dep. on k , crit. pairs

OF OPTIMALS $\leq \binom{k}{k_1} \cdot \left[\underset{\text{values}}{h^{\tau}} \underset{\text{Minimizes}}{(k_1^{\tau} + 1)} \right]^{k_1} \leq \binom{k^{1/2} n^{\epsilon}}{2k_1}$

NOW TAKE

$\tau' :=$ REPLACE IN τ ALL CRITICAL PAIRS IN h_1

BY σ_1

//

τ' DETERMINES $h_2 :=$ the FIRST $h \in H$, $h \parallel \tau'$

//

a gain $\leq (1 + \frac{1}{k} \cdot n^{\frac{1}{k}})^{2k_2}$

options for σ_2

($k_2 := \#$ of CRIT. PAIRS
in h_2)

ALL TOGETHER:

$$\underbrace{(\# \text{ of } \sigma\text{'s})}_{b} \cdot \prod_{i \leq k} (R^{1/2 \cdot u^i \epsilon})^{2u^i} = a (R^{1/2 \cdot u^k \epsilon})^k$$

THIS ROWS THE NO. OF DATA SETS SPECIFIED ANY

$$b := \binom{u+1}{u^{\epsilon - k/2 + 1}} \binom{u}{u^{\epsilon - 2/2}} (u - u^{\epsilon - 2/2})!$$

$$a := \# \text{ of } \rho\text{'s} = \binom{u+1}{u^{\epsilon + 1}} \binom{u}{u^{\epsilon}} (u - u^{\epsilon})!$$

TO REACH A CONTRADICTION FOR ONE A
IT SUFFICES TO VERIFY:

$$a > b \cdot (L^{1/2} \cdot u^3)^4$$

|| _____ FOR ALL $3 \leq j \leq n$

IT SUFFICES IF:

$$a > s \cdot b \cdot (L^{1/2} \cdot u^3)^4$$

ESSENTIAL & **CONDUIT**

FOR THE PARAMETERS, USING $0 < \delta < \epsilon < 1/5$
AND $u \gg 1$.

□
[15:22?]
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