## Surreal numbers

Martin Melicher

March 28, 2021

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- Our order will be different from the order in the book
- We start by defining games and learn to make arithmetics on them
- Then we define surreal numbers as a special kind of games
- Field of surreal numbers turns out to be similar in structure to reals, but much richer

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• If  $a = \{X \mid Y\}$  is a game, we write  $a_L$  for X and  $a_R$  for Y

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#### Definition (Game)

For each ordinal  $\alpha$  we define  $M_{\alpha}$  as follows:

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We say that x is a game, if  $x \in M_{\alpha}$  for some  $\alpha$ . We call the smallest such  $\alpha$  the birthday of x.

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- $|M_2| = 256$
- The class of all games is a proper class.
- If x is a game and  $y \in x_L \cup x_R$ , then birthday(y) < birthday(x).

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There is no infinite sequence of games  $\{g_i\}$  such that  $g_{i+1} \in (g_i)_L \cup (g_i)_R$  for each  $i \in \mathbb{N}$ . If there was, their birthdays would form an infinite decreasing sequence of ordinals.

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#### Lemma (Induction)

Let  $\phi$  be a property such that for each game x:

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$$(\forall y \in x_L \cup x_R : \phi(y)) \rightarrow \phi(x)$$

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#### Lemma (Induction 2)

Let  $\phi$  be a property such that for each x, y:

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$$(\forall x' \in x_L \cup x_R : \phi(x', y)) \land (\forall y' \in y_L \cup y_R : \phi(x, y')) \rightarrow \phi(x, y)$$

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Similar lemma holds for  $\phi(x, y, z)$ , etc.

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# Definition of game

Consequences:

• Every game terminates after finite number of moves.

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Consequences:

- Every game terminates after finite number of moves.
- For given game and given starting player, exactly one player has a winning strategy.

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#### Definition

Let *a* be a game. We say that  $a \ge 0$ , if, given Right starts, Left has a winning strategy. Alternatively,  $a \ge 0$  iff  $\forall b \in a_R \exists c \in b_L : c \ge 0$ .

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Let a, b be games. By a + b we mean a game  $\{(a_L + b) \cup (a + b_L) \mid (a_R + b) \cup (a + b_R)\}$ 

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Let *a* be a game. By -a we mean a game  $\{-a_R \mid -a_L\}$ 

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We define a - b as a + (-b) and  $a \ge b$  as  $a - b \ge 0$ . Furthermore:

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We use  $a \equiv b$ , if a and b are exactly the same  $(a_L = b_L \text{ and } a_R = b_R)$ . Examples:

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Let's prove, that =+==1. That is, 国国=日. We need to cleck two inequalities: ①目用三日: We need 目目田 20, so Right starts, Can left win? 日日日 · 日日日 · 日日日 · 日日日 · 日日日 · 日日日 · 田日日 · 田日日 · (vins) · 田日日 · (vins) 日2日日: We need 日前20 

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#### Lemma

Let a, b, c be games. Then:

 $a+0 \equiv a$ 

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#### Lemma

Let *a*, *b*, *c* be games. Then:

- $\bullet a + 0 \equiv a$
- $a + b \equiv b + a$

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- $\bullet a + 0 \equiv a$
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- $(a+b)+c \equiv a+(b+c)$

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- $\bullet a + 0 \equiv a$
- 2  $a+b\equiv b+a$
- $(a+b)+c \equiv a+(b+c)$
- $-(-a) \equiv a$

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Let *a*, *b*, *c* be games. Then:

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#### Lemma

Let *a*, *b*, *c* be games. Then:

- $\bullet a + 0 \equiv a$
- $a + b \equiv b + a$
- (a+b)+c ≡ a+(b+c)
   -(-a) ≡ a

$$(a+b) \equiv -a-b$$

Before, we defined  $a \ge b$  as  $a - b \ge 0$ . Now we know, that  $a - 0 \equiv a + 0 \equiv a$ , so the definition of  $\ge$  is consistent.

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#### Proof

We prove all of them by induction.

a + 0

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**1** 
$$a + 0 \equiv \{a_L \mid a_R\} + \{ \mid \}$$

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**0** 
$$a + 0 \equiv \{a_L \mid a_R\} + \{\mid\} \equiv \{a_L + 0 \mid a_R + 0\} \equiv \{a_L \mid a_R\} \equiv a_L$$

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$$a+b \\ \equiv \{(a_L+b)\cup(a+b_L) \mid (a_R+b)\cup(a+b_R)\}$$

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We prove all of them by induction.

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② 
$$a + b$$
  
≡ { $(a_L + b) \cup (a + b_L) | (a_R + b) \cup (a + b_R)$ }  
≡ { $(b + a_L) \cup (b_L + a) | (b + a_R) \cup (b_R + a)$ }

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#### Proof

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We prove all of them by induction.

$$a + 0 \equiv \{a_L \mid a_R\} + \{ \mid \} \equiv \{a_L + 0 \mid a_R + 0\} \equiv \{a_L \mid a_R\} \equiv a_L$$

$$= \{(a_L + b) \cup (a + b_L) \mid (a_R + b) \cup (a + b_R)\} \\ = \{(b + a_L) \cup (b_L + a) \mid (b + a_R) \cup (b_R + a)\} = b + a.$$

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So Both sides evaluate to  $(a_L + b + c) \cup (a + b_L + c) \cup (a + b + c_L)$  on the left and analogously on the right.

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$$\begin{array}{l} \mathbf{9} & -(a+b) \\ & \equiv -\{(a_L+b)\cup(a+b_L) \mid (a_R+b)\cup(a+b_R)\} \\ & \equiv \{-(a_R+b)\cup-(a+b_R) \mid -(a_L+b)\cup-(a+b_L)\} \end{array}$$

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#### Lemma

Let *a*, *b* be games. Then:

$$1 a \leq b \text{ iff } -a \geq -b.$$

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#### Lemma

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 iff  $-a \geq -b$ 

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$$a = b$$
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 $a \le b \text{ iff } -a \ge -b.$ 

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$$a = b$$
 iff  $-a = -b$ .

**③**  $a \leq 0$  iff given Left starts, Right can win.

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### Proof

**9** Follows from the definition and  $b - a \equiv -a + b \equiv -a - (-b)$ .

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**③**  $a \leq 0$  iff given Left starts, Right can win.

### Proof

- **9** Follows from the definition and  $b a \equiv -a + b \equiv -a (-b)$ .
- $a \le 0$  is equivalent to  $-a \ge 0$ . Thus in -a, given Right starts, Left can win. But -a is just a with roles switched.

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### Consequence

Let *a* be a game. Then:

• a > 0 iff Left has a winning strategy.

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#### Consequence

Let *a* be a game. Then:

- a > 0 iff Left has a winning strategy.
- a < 0 iff Right has a winning strategy.

#### Consequence

Let *a* be a game. Then:

- a > 0 iff Left has a winning strategy.
- a < 0 iff Right has a winning strategy.
- a = 0 iff second player has a winning strategy.

#### Consequence

Let *a* be a game. Then:

- a > 0 iff Left has a winning strategy.
- a < 0 iff Right has a winning strategy.
- a = 0 iff second player has a winning strategy.
- a || 0 iff first player has a winning strategy.

#### Consequence

Let *a* be a game. Then:

- a > 0 iff Left has a winning strategy.
- a < 0 iff Right has a winning strategy.
- a = 0 iff second player has a winning strategy.
- a || 0 iff first player has a winning strategy.

### Consequence

 $-1 < 0 < 1; * || 0; \{* | *\} = 0$ 

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#### Lemma

Let *a*, *b* be games. Then:

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From point 3 it follows that a = a.

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Thus  $a = b \leftrightarrow a + c = b + c$  and  $(a = b \land b = c) \rightarrow a = c$ .

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#### Definition

Game *a* is a surreal number, if  $a_L$  and  $a_R$  are sets of surreal numbers and for each  $x \in a_L$ ,  $y \in a_R$  we have x < y.

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Class No of all number is a proper class.

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Theorem

Let a be a number,  $b \in a_L$  and  $c \in a_R$ . Then b < a < c.

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If a, b are numbers, then so are a + b and -a.

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#### Theorem

These are all the numbers with finite birthday. +, - and  $\leq$  work as you would expect.

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• This is where we start getting infinite sets.

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• 
$$\left\{\frac{1}{4}, \frac{1}{4} + \frac{1}{16}, \frac{1}{4} + \frac{1}{16} + \frac{1}{64}, \dots \mid \dots, \frac{1}{4} + \frac{1}{16} + \frac{1}{32}, \frac{1}{4} + \frac{1}{8}, \frac{1}{2}\right\} = \frac{1}{3}$$
. Indeed,  
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- This is where we start getting infinite sets.
- $\left\{\frac{1}{4}, \frac{1}{4} + \frac{1}{16}, \frac{1}{4} + \frac{1}{16} + \frac{1}{64}, \dots \mid \dots, \frac{1}{4} + \frac{1}{16} + \frac{1}{32}, \frac{1}{4} + \frac{1}{8}, \frac{1}{2}\right\} = \frac{1}{3}$ . Indeed,  $\frac{1}{3} + \frac{1}{3} + \frac{1}{3} = 1$ .
- All the real numbers.
- {0,1,2,3,... | } =  $\omega$ • { | ...,-3,-2,-1,0} =  $-\omega$ • {0 | ..., $\frac{1}{16}, \frac{1}{8}, \frac{1}{4}, \frac{1}{2}, 1$ } =  $\frac{1}{\omega}$
- $x \pm \frac{1}{\omega}$  for each dyadic fraction x.

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 for each real number  $x$ .  
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- $\{0, 1, 2, ..., \omega \mid \} = \omega + 1$
- $\{0,1,2,\dots ~|~ \omega\} = \omega 1$
- $\bullet$  Analogously  $-\omega-1$  and  $-\omega+1$

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$$\{0,1,2,...\mid \omega-1\}=\omega-2$$
, similarly  $\omega-3$ ,  $\omega-4$ , ...

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- $\{0, 1, 2, \dots \mid \omega 1\} = \omega 2$ , similarly  $\omega 3$ ,  $\omega 4$ , ...
- $\{0, 1, 2, 3, \dots \mid \dots, \omega 3, \omega 2, \omega 1, \omega\} = \frac{\omega}{2}$

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- $\left\{0, 1, 2, 3, \dots \mid ..., \frac{\omega}{8}, \frac{\omega}{4}, \frac{\omega}{2}, \omega\right\} = \sqrt{\omega}$

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• {0,1,2,... | 
$$\omega$$
 - 1} =  $\omega$  - 2, similarly  $\omega$  - 3,  $\omega$  - 4, ...  
• {0,1,2,3,... | ..., $\omega$  - 3, $\omega$  - 2, $\omega$  - 1, $\omega$ } =  $\frac{\omega}{2}$   
• {0,1,2,3,... | ..., $\frac{\omega}{8}, \frac{\omega}{4}, \frac{\omega}{2}, \omega$ } =  $\sqrt{\omega}$   
• {0 | ..., $\frac{1}{8\omega}, \frac{1}{4\omega}, \frac{1}{2\omega}, \frac{1}{\omega}$ } =  $\frac{1}{\omega^2}$ 

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$$\{0, 1, 2, \dots \mid \omega - 1\} = \omega - 2$$
, similarly  $\omega - 3, \omega - 4, \dots$   
•  $\{0, 1, 2, 3, \dots \mid \dots, \omega - 3, \omega - 2, \omega - 1, \omega\} = \frac{\omega}{2}$   
•  $\{0, 1, 2, 3, \dots \mid \dots, \frac{\omega}{8}, \frac{\omega}{4}, \frac{\omega}{2}, \omega\} = \sqrt{\omega}$   
•  $\{0 \mid \dots, \frac{1}{8\omega}, \frac{1}{4\omega}, \frac{1}{2\omega}, \frac{1}{\omega}\} = \frac{1}{\omega^2}$   
•  $\sqrt[3]{\omega + 1} - \frac{\pi}{\omega}, \omega^{\frac{1}{\omega}}, \dots$ 

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#### Simplicity theorem

Let x be a game and z be a number such that all of  $x_L \not\geq z \not\geq x_R$  hold and no element of  $z_L \cup z_R$  satisfies the same condition. Then x = z.

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#### Consequence

Let  $\{L \mid R\}$  be a number of unknown value. Then this number is equal to simplest such x, that all of L < x < R hold. Here 'simplest' means 'with the smallest birthday'.

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