

Getting an ordinary 2-valued structures

(1) General situation:

L : a countable lang.

B : a complete Bool. alg.

A : a B -valued L -structure

H : a L -sentence of special interest
(e.g. the CH in our case).

U.I.O.g. we assume that all parameters
from A occurring in H have names in L .

(2) For each L -f'n $A(x_1, \dots, x_n)$ introduce
a f. symbol $\text{val}_A(x_1, \dots, x_n)$ and interpret
it by the map:

$$\text{val}_A : \vec{f} \in A^n \rightarrow [A(\vec{f})] \in B$$

Note that there are countably many of these
new symbols.

(3) Define a "combined structure" \mathcal{C} :

$$(A, B, \text{val}_A, \dots)$$

(2)

with one sort for A , another for B .
 \mathcal{C} is a first-order ordering (i.e. 2-valued) structure.

Note that the lang. of \mathcal{C} (i.e. \mathcal{L} , the lang. of B plus all vol_A) is countable.

(4) Apply to \mathcal{C} the Löwenheim-Skolem theorem to get its countable elementary substructure $\mathcal{C}' \preceq \mathcal{C}$:

$$\mathcal{C}' : (A', B', \text{vol}_A, \dots)$$

Claim: The truth evaluation of sentences over A' with values in B' defined by:

$$\llbracket A(\vec{x}) \rrbracket' := \text{vol}_A(\llbracket A(\vec{x}) \rrbracket)$$

for $\vec{x}_i \in A'$, satisfies all conditions required.

Proof: Let us check e.g. the condition about \exists . We want:

$$\llbracket \exists y B(\vec{x}, y) \rrbracket' = \sup_{y \in A'} \llbracket B(\vec{x}, y) \rrbracket'$$

This can be expressed using the language with val - symbols by the sentences:

$$\forall x \forall y, \text{val}_{\exists y B(x,y)}(x) \geq \text{val}_{B(x,y)}(x,y)$$

and

$$\forall b \forall x (\text{val}_{\exists y B(x,y)}(x) > b \rightarrow \exists y \text{val}_{B(x,y)}(x,y) > b)$$

where x, y range over A' -set and b over B' -set.

Both these sentences are true in \mathcal{F} and hence also in \mathcal{F}' .

□ done.

(5) We shall need the (corollary of the) Rasiowa-Sikorski thm: Let \mathcal{B} be a complete Bool. algebra, $\mathcal{U} \subseteq \mathcal{P}(\mathcal{B})$ a countable collection of its subsets.

Then there exists a homomorphism

$$h : \mathcal{B} \rightarrow \{0,1\} \text{ (= the 2-element obj.)}$$

s.t. for all $u \in \mathcal{U}$:

$$h(\sup u) = \sup \{h(u) \mid u \in \mathcal{U}\}.$$

(6) We shall apply it to \mathbb{B}' and to the following collection $\mathcal{U} \subseteq \mathcal{P}(\mathbb{B}')$:

- for each L -fca $B(x, y)$ and each n -tuple ξ of parameters for A' include in \mathcal{U} the set

$$U_{B(\xi, y)} := \{ \bigvee B(\xi, z) \mid z \in A' \}$$

Note that ~~the set \mathcal{U} is not a σ -algebra~~

$$\sup U_{B(\xi, y)} = \bigvee \exists y B(\xi, y)$$

Let $h: \mathbb{B}' \rightarrow \{0, 1\}$ be a homomorphism guaranteed to exist by the R.-S. thm.

(7) Define an ordinary (2-valued)

L -structure A^* by:

- factor A' by the equiv.-rel. \approx :

$$\xi_1 \approx \xi_2 \iff h(\bigvee \xi_1 = \xi_2) = 1$$

- define for each atomic L-sentence ϕ with parameter ~~from~~ \mathcal{M}^* :

$$\mathcal{M}^* \models \phi \iff \text{of. } h(\prod A \prod') = 1$$

Remark: Elements of \mathcal{M}^* are blocks of \mathcal{M} and we represent them in the right-hand-side by any representative of the respective blocks.

(8) Claim: For any L-sentence B with parameters for \mathcal{M}^* : $\mathcal{M}^* \models B \iff h(\prod B \prod') = 1$.

In particular,

(i) If $\prod B \prod' = 1_{\mathcal{M}}$ in \mathcal{M} then $\mathcal{M}^* \models B$.

(It is a consequence of the theory \mathcal{M} -valid in \mathcal{M} is true in \mathcal{M}^* .)

(ii) $\mathcal{M}^* \models \neg H$.

(In particular, CH fails in \mathcal{M}^* in our case.)

Proof: (ii) is proved by induction on the logical complexity of B . The cases of \neg and \wedge need (*) of (5) for all $\mathcal{M}(\xi_i, \gamma)$ for (6).

(i): By (1), $\prod H \prod' = 0_{\mathcal{M}}$, so (i) implies $\mathcal{M}^* \models \neg H$.

□