Using Automated Theorem Provers in Non-associative Algebra

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LPAR 2008, Doha, Qatar

[The authors] demonstrate that (contrary to the view amongst some in AR), provided a sufficiently effective AR tool is available, there are some mathematicians who will indeed use such a tool.

- anonymous referee at ESARM

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This talk

- is about *solving open problems* by (first order) automated theorem provers
- is not about formal verification or theory formation, no toy examples

Areas of algebra

- simple axiomatization projects (about 10 papers, since early 90's)
- lattices with operators (about 10?)
 - Robbins problem
 - algebraic logic
- non-associative algebra
 - quasigroups and loops (about 25, since 1996)
 - etc. (about 5)

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Types of computation

- direct proofs of difficult open problems
- proving tedious technical steps
- quickly checking easy conjectures
- exhaustive search

Main problems

formalization in FOL

- almost nothing formalizable directly
- sometimes a highly non-trivial task
- which formalization is optimal

finding a proof

- which prover, setting up parameters
- reading and understanding the proof
 - yes, we want to understand it! (usually)
 - simplifying the proof
 - improving readability (introducing concepts, lemmas, etc.)

For every distribuive groupoid G, there exists a congruence α of G such that G/α is symmetric and all blocks of α are medial.

cnf(sos,axiom,mult(A,mult(B,C)) = mult(mult(A,B),mult(A,C))).
cnf(sos,axiom,mult(mult(A,B),C) = mult(mult(A,C),mult(B,C))).

cnf(goals,negate_conjecture,mult(mult(mult(a,b),mult(c,d)), mult(mult(a,c),mult(b,d))) != mult(mult(mult(a,c),mult(b,d)), mult(mult(a,b),mult(c,d)))).

cnf(goals,negated_conjecture,mult(mult(mult(a,b),mult(c,d)), mult(mult(mult(a,b),mult(c,d)),mult(mult(a,c),mult(b,d)))) != mult(mult(a,c),mult(b,d))). Bruck loops with abelian inner mapping group are centrally nilpotent.

cnf(sos,axiom,mult(unit,A) = A).cnf(sos,axiom,mult(A,unit) = A).cnf(sos,axiom,mult(A,i(A)) = unit). cnf(sos,axiom,mult(i(A),A) = unit).cnf(sos,axiom,i(mult(A,B)) = mult(i(A),i(B))). cnf(sos,axiom,mult(i(A),mult(A,B)) = B).cnf(sos,axiom,rd(mult(A,B),B) = A).cnf(sos,axiom,mult(rd(A,B),B) = A).cnf(sos,axiom,mult(mult(A,mult(B,A)),C) = mult(A,mult(B,mult(A,C)))). cnf(sos,axiom.mult(mult(A,B),C) =mult(mult(A,mult(B,C)),asoc(A,B,C))). $cnf(sos,axiom,op_l(A,B,C) =$ mult(i(mult(C,B)),mult(C,mult(B,A)))). cnf(sos,axiom,op_r(A,B,C) = rd(mult(mult(A,B),C),mult(B,C))). cnf(sos,axiom,op_t(A,B) = mult(i(B),mult(A,B))). $cnf(sos,axiom,op_r(op_r(A,B,C),D,E) = op_r(op_r(A,D,E),B,C)).$ $cnf(sos,axiom,op_l(op_r(A,B,C),D,E) = op_r(op_l(A,D,E),B,C)).$ $cnf(sos,axiom,op_l(op_l(A,B,C),D,E) = op_l(op_l(A,D,E),B,C)).$ $cnf(sos,axiom,op_t(op_r(A,B,C),D) = op_r(op_t(A,D),B,C)).$ $cnf(sos,axiom,op_t(op_l(A,B,C),D) = op_l(op_t(A,D),B,C)).$ $cnf(sos,axiom,op_t(op_t(A,B),C) = op_t(op_t(A,C),B)).$

cnf(goals,negated_conjecture,asoc(asoc(a,b,c),d,e) != unit).

Our work, so far

- proving theorems
- QPTP library

$\mathsf{QPTP} = \mathsf{Quasigroup}$ problems for theorem provers

- = a collection of results in loop theory obtained with assistance of ATP
- all papers covered, about 100 problems selected (about 80% equational)
- both formal (TPTP) and informal (paper) description
- o downloadable at www.karlin.mff.cuni.cz/~stanovsk/qptp
- a benchmark (selected provers from CASC): Waldmeister >> E, Gandalf, Prover9, Vampire >> Spass

Summary

- (some) mathematicians use automated theorem provers
- ATPs can prove difficult theorems
- If you have a software that could solve our problems, let me know immediately!