

Some consequences: (in an arbitrary semigroup)

(8)

- \mathcal{H} -classes contain ≤ 1 idempotent
- maximal subgroups are precisely \mathcal{H} -classes w/ idempotent
- max. subgroups are disjoint
- S idempotent \rightarrow \mathcal{H} -classes have 1 elt., i.e., $\mathcal{H} = id$
- S group $\rightarrow \mathcal{H} = G \times G$

Regularity

a regular \equiv ~~the~~ the following equivalent conditions hold:

- (1) $\exists b \quad aba = a$
- (2) $\exists a' \quad a'a = a \quad \& \quad a'a' = a'$
- (3) $L_a \cap E_S \neq \emptyset$
- (4) $R_a \cap E_S \neq \emptyset$

Proof: (1) \Rightarrow (2) $\hat{a} := bab$

(2) \Rightarrow (1) trivial

(1) \Rightarrow (3) $e := ba$

(3) \Rightarrow (1) $e \in L_a \cap E_S \Rightarrow \exists b \quad ba = e \Rightarrow aba = ae = a$

(4) dually

$Reg(S) := \{a \in S : a \text{ regular}\}$

S regular $\equiv Reg(S) = S$

Ex.:

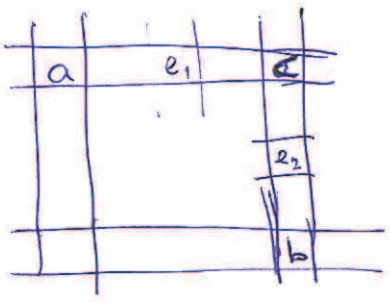
- idempotent semigroups
- groups
- T_X (see also eggbox for T_3)

(*) NOTE: all (1)-(4) hold for a nice reason

~~a regular $\Leftrightarrow \exists b \{aba = a\}$ associates of a
 $\forall a \in S \{a' \{a'a = a \ \& \ a'a' = a'\}$ inverses of a~~

Prop.: $Reg(S)$ is union of \mathcal{D} -classes, i.e., $\left. \begin{matrix} a \in Reg(S) \\ a \mathcal{D} b \end{matrix} \right\} \Rightarrow b \in Reg(S)$

Proof:



$$a \mathcal{D} b \stackrel{\text{Proj.}}{\Leftrightarrow} \exists c \ aRc \ \& \ b$$

$a \text{ regular} \Rightarrow \exists e_1 \in R_a \cap E_S = R_c \cap E_S$
 $\Rightarrow c \text{ regular}$
 $\Rightarrow \exists e_2 \in R_c \cap E_S = L_b \cap E_S$
 $\Rightarrow b \text{ regular} \quad \square$

- iii) $S \text{ regular} \Rightarrow S/\alpha \text{ regular}$
- iii) $S_i \text{ regular} \Rightarrow \prod S_i \text{ regular}$
- iii) $S \text{ regular}, T \leq S \not\Rightarrow T \text{ regular}$
 $\dots (\mathbb{N}, +) \leq (\mathbb{Z}, +)$

$As(a) := \{ b : aba = a \} \quad \dots \text{associates of } a$
 $V(a) := \{ a' : aa'a = a, a'a'a = a' \} \quad \dots \text{inverses of } a$

iii) $V(a) \subseteq As(a), \quad As(a) \rightarrow V(a)$
 $b \mapsto bab$

Lemma:

$As(a) \rightarrow E_S \cap L_a$ is onto
 $b \mapsto ba$

Proof: $e \in E_S \cap L_a \Rightarrow \left. \begin{array}{l} ae = a \\ \exists x \ xa = e \\ b := xax \end{array} \right\}$

~~$V(a) \subseteq E_S \cap L_a$ is onto~~
 ~~$b \mapsto ba$~~
 ~~$a \in V(a) \Rightarrow a = xax$~~

~~$a \in As(a) \Rightarrow a = xax$~~
 $ba = xaxa = ee = e \quad \checkmark$
 $? b \in V(a)?$
 $aba = axaxa = ae^2 = ae = a$
 $bab = xaxaxa = e^3x = ex = xax = a \quad \square$

HENCE: $|E_S \cap L_a| \leq |V(a)| \leq |As(a)|$ (a regular = all $\neq 0$)

Prop.: a regular $\Rightarrow \left[\begin{array}{l} |V(a)|=1 \Leftrightarrow \left\{ \begin{array}{l} |E_s \cap L_a|=1 \\ |E_s \cap R_a|=1 \end{array} \right. \end{array} \right] \quad (10)$

Proof: \Rightarrow from $\textcircled{10}$ above

$$\left. \begin{array}{l} \textcircled{\Leftarrow} \text{ let } a', a'' \in V(a) \\ E_s \cap L_a = \{e\} \\ E_s \cap R_a = \{f\} \end{array} \right\} \begin{array}{l} a'a = a''a = e \\ aa' = aa'' = f \end{array} \quad \left(\begin{array}{l} \text{by } V(a) \rightarrow \dots \\ \text{(by } b \mapsto ba) \end{array} \right)$$

\Downarrow

$$a = \underline{a'a} = a''a' = a''a'' = a'' \quad \square$$

Prop.: a, b regular \Rightarrow

~~$a \preceq b \Leftrightarrow \forall a' \in V(a) \exists b' \in V(b) \text{ s.t. } a'a = b'b$~~

$$\begin{aligned} a \preceq b &\Leftrightarrow (1) \forall a' \in V(a) \exists b' \in V(b) \quad a'a = b'b \\ &\Leftrightarrow (2) \exists a' \in V(a) \exists b' \in V(b) \quad a'a = b'b \end{aligned}$$

Proof: ~~the following is not correct:~~

$$a \preceq b \Rightarrow (1): \quad \left. \begin{array}{l} a = xb \\ b = ya \end{array} \right\} \begin{array}{l} a'a = \underline{a'xb} \\ b' := a'x \quad ? b' \in V(b)? \end{array}$$

$$b'b' = \underline{b'a'xb} = ya'a = ya = b$$

$$\underline{b'b'} = \underline{a'xb'a'x} = a'a'a'x = a'x = b' \quad \blacksquare$$

(1) \Rightarrow (2) trivial

$$\begin{aligned} (2) \Rightarrow \textcircled{1} : a \preceq b : a'a = b'b &\Rightarrow \begin{array}{l} a = a'a = \underline{ab'b} \\ \underline{ba'a} = b'b'b = b \end{array} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} a \preceq b \quad \square \end{aligned}$$

Inverse semigroups

$$\equiv \forall a \exists ! a' \quad a a' a = a \quad \& \quad a' a a' = a'$$

i.e., $|V(a)| = 1 \quad \forall a$

~~Ex:~~ Ex: \circ groups, semilattices
 $\circ I_X =$ partial bijections $= \{ f \in PT_X : f \text{ 1-1} \} \subseteq PT_X$

Note: $(a')' = a, (ab)' = b'a$... (not easy to prove!! (later))
 " — either!! (later)

Theorem: TFAE: (1) ~~inverse~~ inverse
 (2) every \mathcal{L} -class contains exactly one idempotent
 & every \mathcal{R} -class " "
 (3) regular & all idempotents commute
 (\circledast hence $E_S \leq S$)
 ... $ef \cdot ef \stackrel{\text{comm.}}{=} ee \cdot ff = ef$

Proof: (1) \Leftrightarrow (2) see Prop. on (10)

(3) \Rightarrow (2) at least 1 ... from regularity
 at most 1 ... $e \neq f \Rightarrow e = ef = fe = f$
 (\uparrow Zermod)

(2) \Rightarrow (3) $e, f \in E_S \quad z := (ef)'$
 (the unique inverse)

Claim 1: $fze = z$

... $(ef)(fze)(ef) = \underline{efzef} = ef$
 $(fze)(ef)(fze) = \underline{fzefze} = fze$ } fze is inverse to ef
 $\Rightarrow fze = z$

Claim 2: $ef \in E_S$

... ~~zz = fze fze = fze = z~~

$\Rightarrow z$ is inverse for z } $ef = z \in E_S$
 ef — " — z

$\circledast u \in E_S \Rightarrow u' = u$

Claim 3: $fe \in E_S$

... dually

$$\Rightarrow \left. \begin{aligned} (ef)(fe)(ef) &= efef = ef \\ (fe)(ef)(fe) &= fefe = fe \end{aligned} \right\} ef = fe \quad \square$$

Basic properties of inverse semigroups:

(1) $(ab)' = b'a'$

(2) $(a')' = a$

(3) $e \in E_S \Rightarrow aea' \in E_S \quad \& \quad E_S \leq S$

(4) $a \not\sim b \Leftrightarrow a'a = b'b$

$a \sim b \Leftrightarrow aa' = bb'$

(5) $e, f \in E_S, e \not\sim f \Rightarrow \exists a \quad a'a = e \quad \& \quad aa' = f$

(in particular, for $e = f \in E_S: \exists a \quad e = a'a = aa'$)

Proof: (1) $\underbrace{(ab)(b'a')(ab)}_{e \in E_S, e \in E_S} \stackrel{\text{comm}}{=} a'a b'b' = ab$ } \Rightarrow unique' $(ab)' = b'a'$

$\underbrace{(b'a')(ab)(b'a')}_{e \in E_S, e \in E_S} \stackrel{\text{comm}}{=} b'a'$

(2) a is inverse for $a' \Rightarrow (a')' = a$

(3) $e, f \in E_S \Rightarrow ef \cdot ef \stackrel{\text{comm}}{=} ee ff = ef$

$\Rightarrow \underbrace{aea'aea'}_{e \in E_S} = \underbrace{aa'aea'}_{e \in E_S} = aea'$

(4) see Prop. on (10)

(5) $\exists a \quad e \not\sim a \not\sim f$

\downarrow (4) \Downarrow (4)

$a'a = ee = e \quad aa' = ff = f$

⊗ Hint: prove that if $e \in E_S$ then $e = ee'$. Hence idempotents commute.

□

→ EQUATIONAL BASE

Coro: $(S, \cdot, ')$ is an inverse semigroup \Leftrightarrow

$\left. \begin{aligned} \bullet x(yz) &= (xy)z \\ \bullet xx'x &= x, \quad x'xx' = x' \\ \bullet (xx')(yy') &= (yy')(xx') \end{aligned} \right\} \forall x, y, z$

Proof: \Rightarrow use Thm (3) & property (5) \in HW1

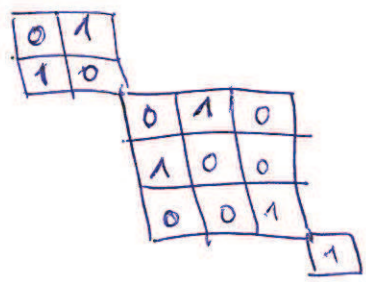
Eggbox for inverse semigroups:

... every row / column exactly one idempotent

→ $\text{put } \leq_1^0$ # of idp. in an \mathcal{H} -class

→ permutation matrices
 \uparrow
 \mathcal{D} -classes =

Ex:



Prop: I_X is an inverse semigroup

Proof: $I_X = \{ \alpha : A \rightarrow B \text{ bijective} : A, B \subseteq X \}$

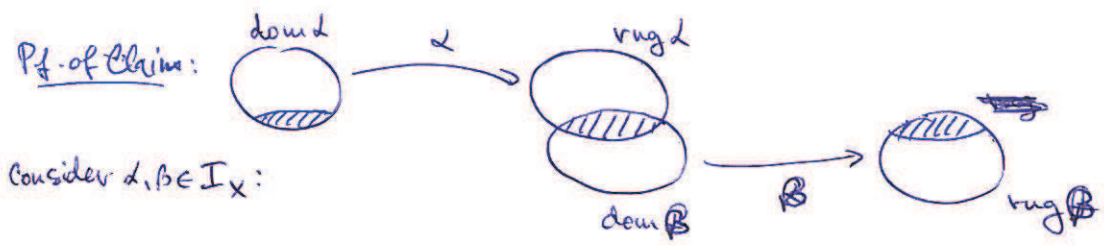
\exists
 $\alpha' := \alpha^{-1} : B \rightarrow A$

we have $\left. \begin{matrix} \alpha \alpha' = id_B \\ \alpha' \alpha = id_A \end{matrix} \right\} \Rightarrow \begin{matrix} \alpha' \alpha \alpha' = \alpha' \\ \alpha \alpha' \alpha = \alpha \end{matrix} \Rightarrow \text{regular}$

... ? unique ? ... difficult
 ... ? idempotents commute ?

Claim: $\alpha \alpha = \alpha \Leftrightarrow \exists A \subseteq X \quad \alpha = id_A$

Then it is trivial: $id_A id_B = id_{A \cap B} = id_B id_A \quad \square$



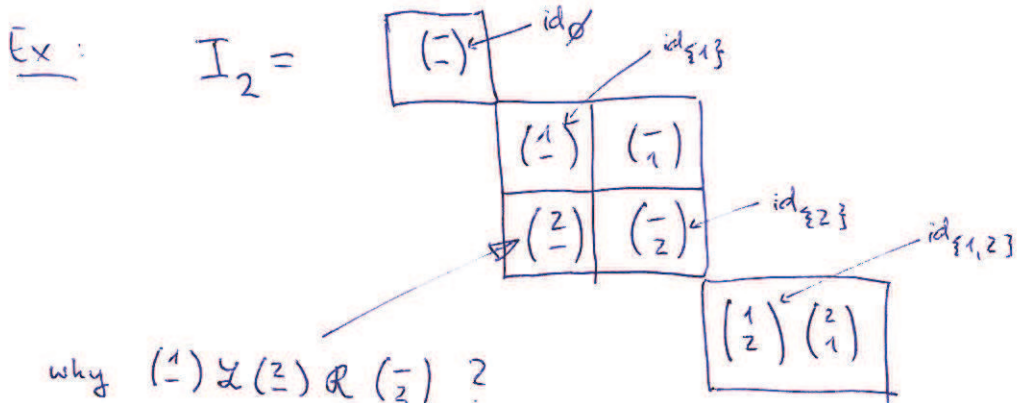
$\text{dom } \beta \alpha = \alpha^{-1}(\text{rng } \alpha \cap \text{dom } \beta)$
 $\text{rng } \beta \alpha = \beta(\text{rng } \alpha \cap \text{dom } \beta)$

\Rightarrow We ~~assume~~ $\alpha = \beta$ & $\alpha \alpha = \alpha$, hence

$\alpha^{-1}(\text{rng } \alpha) = \text{dom } \alpha = \text{dom } \alpha \alpha = \alpha^{-1}(\text{rng } \alpha \cap \text{dom } \alpha)$
 $\alpha(\text{dom } \alpha) = \text{rng } \alpha = \text{rng } \alpha \alpha = \alpha(\text{---})$

$\xRightarrow{\alpha^{-1}}$ $\text{rng } \alpha = \text{rng } \alpha \cap \text{dom } \alpha = \text{dom } \alpha$

$\Rightarrow \alpha$ is permutation on $A \xRightarrow{\alpha \alpha = \alpha} \alpha = id_A \quad \square$



$\mathcal{L}: \begin{pmatrix} 2 \\ - \end{pmatrix} \begin{pmatrix} 2 \\ - \end{pmatrix} = \begin{pmatrix} - \\ 1 \end{pmatrix} \begin{pmatrix} 2 \\ - \end{pmatrix} = \begin{pmatrix} 1 \\ - \end{pmatrix}$

$\mathcal{R}: \begin{pmatrix} 2 \\ - \end{pmatrix} \begin{pmatrix} - \\ 2 \end{pmatrix} = \begin{pmatrix} 2 \\ - \end{pmatrix} \begin{pmatrix} - \\ 1 \end{pmatrix} = \begin{pmatrix} - \\ 2 \end{pmatrix}$

Vagner - Preston theorem: S inverse sug. $\Rightarrow S \hookrightarrow I_S$
 $a \mapsto \lambda_a \uparrow a^s$
 1952 USSR 1954 USA

Ex: groups ... $a^s = S$, hence $S \hookrightarrow$ symmetric group over S
 $a \mapsto \lambda_a$
 Cayley rep.

Ex: \wedge -semilattices ... $a^s = aS = \downarrow a = \{x : x \leq a\}$
 \bullet a idempotent $\Rightarrow \varphi(a)$ idempotent for any φ hom.
 $\Rightarrow \lambda_a \uparrow \downarrow a = id_{\downarrow a}$
 (indeed, $\lambda_a(x) = ax = x \quad \forall x \leq a$)
 $\bullet id_A \circ id_B = id_{A \cap B}$

Hence, in fact, $S \hookrightarrow (\mathcal{P}(S), \wedge)$
 $a \mapsto \downarrow a$... semilattice of subsets

Proof: $\sigma_a := \lambda_a \uparrow a^s$, denote φ the mapping $a \mapsto \sigma_a$

1) $\sigma_a : a^s \rightarrow aS$... $\lambda_a(a^s) = a^s \in aS$
 $\underbrace{\hspace{10em}}_{\text{onto}} \quad \quad \quad aS \ni as = aa^s = \lambda_a(a^s)$
 $\hspace{15em} \underbrace{\hspace{10em}}_{\in a^s}$

2) $\sigma_a \in I_S$... $\sigma_a(x) = \sigma_a(y)$ for some $x, y \in a^s$
 (i.e., it is 1-1) $\underbrace{\hspace{10em}}_{x=a^s, y=a^t}$

$$\Rightarrow aas = aat$$

$$\Rightarrow daas = daat$$

$x = a's \quad \quad \quad a't = y$

3) φ is 1-1 : $\sigma_a = \sigma_b \Rightarrow \text{dom } \sigma_a = \text{dom } \sigma_b = b's$

$$\Rightarrow a' \mathcal{R} b'$$

Prop. (4)

$$\Leftrightarrow a'a'' = b'b'' = b'b \in a's = b's$$

$$\Rightarrow \sigma_a(a'a) = \sigma_b(b'b)$$

$a = a'a \quad \quad \quad b'b'b = b$

4) φ is hom. : ? $\sigma_a \circ \sigma_b = \sigma_{a \cdot b}$?

4a) Indeed $\sigma_a \circ \sigma_b = \sigma_{a \cdot b}$, since $\sigma_a(\sigma_b(x)) = a(bx)$
 $\forall x \in S \quad \sigma_{a \cdot b}(x) = (ab)x$ \gg

4b) $\text{dom } \sigma_a \circ \sigma_b \stackrel{?}{=} \text{dom } \sigma_{a \cdot b} = (ab)'S$?

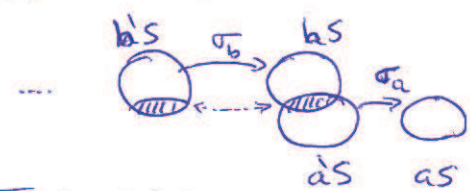
$\text{dom } \sigma_a \circ \sigma_b = \sigma_b^{-1}(b's \cap a's)$

$= \sigma_b^{-1}(bb's \cap a'a's)$

$\stackrel{\textcircled{*}}{=} \sigma_b^{-1}((bb')(a'a)S)$

$\stackrel{\textcircled{\Delta}}{=} \sigma_b^{-1}((bb')(a'a)S)$

$= b'bb'a'aS = b'a'aS = b'a'S = (ab)'S \checkmark$



$b's = bb's$
 because $b = bb'b$
 $a's = a'a's$
 because $a = a'aa'$

$\textcircled{*} \quad eS \cap fS = efS \quad \forall e, f \in E_S :$

$\textcircled{\ominus} \quad efS \subseteq eS, feS \subseteq fS$

$\textcircled{\subseteq} \quad z \in eS \cap fS$

$\Rightarrow z = ex = fy$

$\Rightarrow efz = efly = efy = ez = eex = ex = z$

$\Rightarrow z \in efS$

$\textcircled{\Delta} \quad ? \quad \sigma_b \sigma_b^{-1} = id_{b's}, \sigma_b^{-1} \sigma_b = id_{a's} ?$

$\sigma_b^{-1} : b'S = bS \rightarrow b'S$

\Rightarrow dom, rng OK

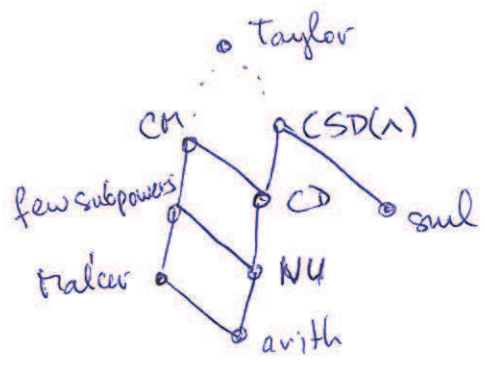
$\sigma_b \sigma_b^{-1}(bs) = bb'bs = bs \checkmark$

$\sigma_b^{-1} \sigma_b(b's) = b'bb's = b's \checkmark$

□

Inverse semigroups & Maltsev conditions

Ex: groups - are ~~not~~ ^{Maltsev} Maltsev, no sul-term, not CSD(n) HW
 semilattices - sul-term, not CM HW



Taylor term: idemp. & $t\left(\begin{smallmatrix} x & & \\ & x & \\ & & x \end{smallmatrix}\right) = t\left(\begin{smallmatrix} & y & \\ & & y \\ x & & \end{smallmatrix}\right)$

$$t(xyzuv) := xy'zu'v$$

... idemp. : $\underline{x}x'x'x'x = x'x'x = x$

.. $t(xxyy\underline{x}) = t(yyxx\underline{x})$

$t(\underline{x}xyy) = t(\underline{y}yyxx)$

... ~~the~~ coord. 1..4
 ... coord. 3..5

Congruences of inverse semigroups

\leftrightarrow (normal subsemigroup, normal congruence on E_S)

α congruence on $S \xrightarrow{\text{inverse}} \text{trace } \text{tr } \alpha = \alpha \cap E_S^2$

kernel $\text{Ker } \alpha = \bigcup_{e \in E_S} [e]_\alpha$

\Rightarrow Ex.: groups: $\text{tr } \alpha$ triv., $\text{Ker } \alpha = \{1\}$ // semilattices: $\text{tr } \alpha = \alpha$, $\text{Ker } \alpha = E_S$

def.: $T \leq S$ is called full $\equiv E_T = E_S$

self-conjugate $\equiv \forall t \in T \forall a \in S \text{ at } a' \in T$

normal \equiv full, self-conjugate, inverse

Lemma:

~~Kernel~~ $\text{Ker } \alpha$ is normal

Pf: full \checkmark

$\text{Ker } \alpha \leq S$?

$a \in [e], b \in [f] \rightarrow abc \in [ef]$

$ade \quad bdf \quad abdef \in E_S$

~~Kernel~~ self-conjugate ?

$a \in [e], b \in S \Rightarrow bab' \in [beb']$

$ade \quad bab' \alpha beb' \in E_S$

inverse ?

$a \in [e] \Rightarrow a' \in [e]$
 $ade \quad a'de' = e$

\square

def.: $T \leq S$ normal, $\sigma \in \text{Con } T$ is called normal $\equiv \forall a, b \in T \forall c \in S$

$a \sigma b \Rightarrow ca'c' \sigma cbc'$

Lemma: $\text{tr } \alpha$ is normal (on $E_S \leq S$ normal)

Pf: $a(\text{tr } \alpha) b \Rightarrow a \alpha b \Rightarrow ca'c' \alpha cbc' \Rightarrow ca'c'(\text{tr } \alpha) cbc'$
 $a, b \in E_S \quad ca'c', cbc' \in E_S$

\square

Ex.: S/α is a group iff $\text{tr } \alpha = E_S^2$

Proof: ~~Kernel~~

\Rightarrow obvious

\Leftarrow from Lallement

Note: S inverse $\Rightarrow \left\{ \begin{array}{l} S \text{ group} \Leftrightarrow |E_S| = 1 \\ \Leftarrow \dots \text{ look at eggbox} \end{array} \right.$

\square

Theorem: S inverse smg.

(1) $\alpha \in \text{Con } S \Rightarrow (\text{Ker } \alpha, \text{tr } \alpha)$ is a congruence pair

(2) (N, τ) congruence pair \Rightarrow

$$\alpha_{(N, \tau)} := \{ (a, b) : (a'a, b'b) \in \tau \ \& \ ab' \in N \} \in \text{Con } S$$

(3) $\text{Ker } \alpha_{(N, \tau)} = N, \text{tr } \alpha_{(N, \tau)} = \tau$

$$\alpha_{(\text{Ker } \alpha, \text{tr } \alpha)} = \alpha$$

def.: (N, τ) congruence pair \equiv

- $N \leq S$ normal
- $\tau \in \text{Con}(E_S)$ normal

$$\left. \begin{array}{l} \text{(CP1)} \ ae \in N \\ \quad (e, a'a) \in \tau \end{array} \right\} \Rightarrow ae \in N$$

$$\text{(CP2)} \ ae \in N \Rightarrow (aa', a'a) \in \tau$$

Ex.: in groups: $N \leq S, \tau$ trivial on E_S trivial

$$\alpha \rightsquigarrow N = [1], N \rightsquigarrow \alpha = \{ (a, b) : \underbrace{a'a = b'b}_{\text{triv.}} \ \& \ ab' \in N \}$$

in semilattices:

$$N \leq S \text{ full} \rightsquigarrow N = S, \alpha = \{ (a, b) : (a, b) \in \tau \ \& \ \underbrace{ab \in N}_{\text{triv.}} \} \text{ hence } \alpha = \tau$$

Proof: (1) $\text{Ker } \alpha \leq S$ normal, $\text{tr } \alpha \in \text{Con}(E_S)$ normal ... see Lemmas

$$\left. \begin{array}{l} \text{(CP1)} \ ae \in [f] \in \text{Ker } \alpha \\ \quad e \in \tau \ a'a \end{array} \right\} \Rightarrow a = a(a'a) \ \& \ ae \ \& \ f \Rightarrow a \in [f] \subseteq \text{Ker } \alpha$$

$$\begin{aligned} \text{(CP2)} \ ae \in [f] \in \text{Ker } \alpha &\Rightarrow aa' \ \& \ ff = f \\ &\quad a \ \& \ f \quad \quad \quad aa' \ \& \ ff = f \\ &\Rightarrow aa' \ \& \ a'a \end{aligned}$$

(2) reflex.: $aa' \in N$ because full

sym.: $ab' \in N \Rightarrow (ab')' \in N$ because inverse
 \parallel
 ba'

trans.: $ab' \in N, bc' \in N \Rightarrow \underbrace{abb'}_{\in E_S} c' \in N \stackrel{\otimes}{\Rightarrow} ac' \in N$

⊗ ≡ $aeb \in N, e \tau a \Rightarrow ab \in N$

$$\left. \begin{aligned} \dots aeb &= \underline{aebb'b} = ab\underline{b'eb} = abf \\ & \qquad \qquad \qquad =: f \in E_S \end{aligned} \right\} \text{(CP1)} \Rightarrow ab \in N$$

$$f = \underline{b'eb} \tau \underline{b'aab} = (ab)'ab$$

↑
because τ is normal

invariant: assume $a \tau b$, prove $ca \tau cb$
 $ac \tau bc$
 $a \tau b$
 $ab \in N$

ca τ cb: $(ca)'(ca) = \underline{a'ca} \tau \underline{b'cb} = (cb)'(cb)$
 (⊗)

$(ca)(cb)' = \underline{cab'c'} \in N$ because self-conj.

ac τ bc: $(ac)'(ac) = \underline{c'ac} \tau \underline{c'bbc} = (bc)'(bc)$
 ↑
 because τ normal

$(ac)(bc)' = \underline{acc'b} = \underline{acc'b^*b}$
 $= \underline{ab^*c} \underline{c'b} \in N$
 $\in N \quad \in E_S \subseteq N$
 because full

⊕ ≡ $a \tau b, ab \in N \Rightarrow a \tau ea \tau beb \quad \forall e \in E_S$

$$\begin{aligned} \dots a \tau ea &= a \tau ea (a \tau a) \tau ea && \dots \text{idemp.} \\ &\tau \underline{a \tau ea} (b \tau b) \tau \underline{a \tau ea} && \dots \tau \text{ normal} \\ &= a \tau e ab' (ab')' \tau ea \\ &\tau \underline{a \tau e} (ab')' \tau \underline{ab' \tau ea} && \dots \text{by (CP2) \& } \tau \text{ normal} \\ &= \underline{a \tau eb} \tau \underline{a \tau ab'} \tau ea \\ &= a \tau (ab'e)' (ab'e) \tau a && \text{NOTE: } ab'e \in N \text{ since full} \\ &\tau a \tau (ab'e) (ab'e)' \tau a && \dots \text{by (CP2) \& } \tau \text{ normal} \\ &= \underline{a \tau a} \tau \underline{b' \tau ee} \tau \underline{b \tau a \tau a} \\ &\tau \underline{b' \tau b} \tau \underline{b' \tau e} \tau \underline{b \tau b' \tau b} = b' \tau eb \quad \checkmark \end{aligned}$$

(3) $\text{Ker } \alpha_{(N, \tau)} = N$:

(\subseteq) $a \in \text{Ker } \alpha \Rightarrow \begin{matrix} a \in [e] \\ (a \neq 1) \end{matrix} \stackrel{\text{dit}}{\Rightarrow} \left. \begin{matrix} a \cdot a \neq e = e \\ a e' = a e \in N \end{matrix} \right\} \begin{matrix} (\text{ep1}) \\ \Rightarrow \end{matrix} a \in N$

(\supseteq) $\left. \begin{matrix} a \in N \\ e = a \cdot a \end{matrix} \right\} \Rightarrow \text{? } a \in e \text{ ? } \text{ hence } a \in \text{Ker } \alpha$

$\left. \begin{matrix} \uparrow \\ \& \end{matrix} \right\} \begin{matrix} a \cdot a = e = e \cdot e, \text{ hence } a \cdot a \tau e \\ a e' = a \cdot a = a \in N \end{matrix}$

$\text{tr } \alpha_{(N, \tau)} = \tau$:

$\alpha_{(\text{Ker } \alpha, \text{tr } \alpha)} = \alpha$:

} HW

□